# Markovian Models for Genome Rearrangement Evolution 

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## Outline

- Genome Rearrangement Evolution
- The GNT Model
- Distribution of evolutionary distances
- Breakpoint distance
- Inversion distance
- Simulation study: accuracy of tree reconstruction
- Future work


## Genomes As Signed Permutations



## Genomes Evolve by Rearrangements


$\begin{array}{llllllllll}1 & 2 & -6 & -5 & -4 & -3 & 7 & 8 & 9 & 10\end{array}$

Transposition:
$\begin{array}{llllllllll}1 & 2 & 7 & 8 & 3 & 4 & 5 & 6 & 9 & 10\end{array}$

Inverted Transposition:
$\begin{array}{llllllllll}1 & 2 & 7 & 8 & -6 & -5 & -4 & -3 & 9 & 10\end{array}$

## Our Model: the Generalized Nadeau-Taylor Model [STOC'01]

- Three types of events:
- Inversions (INV)
- Transpositions (TRP)
- Inverted Transpositions (ITP)
- Events of the same type are equiprobable
- Probabilities of the three types have fixed ratio

$$
\begin{aligned}
& \operatorname{Pr}(r \in I N V): \operatorname{Pr}(r \in T R P): \operatorname{Pr}(r \in I T P) \\
= & (1-\alpha-\beta): \alpha: \beta
\end{aligned}
$$

- We focus on signed circular genomes in this talk.


## Edit Distances Between Genomes

- (I NV) Inversion distance [Hannenhalli \& Pevzner 1995]
- Computable in linear time [Moret et al 2001]
- (BP) Breakpoint distance [Watterson et al. 1982]
- Computable in linear time
- NJ (BP): [Blanchette, Kunisawa, Sankoff, 1999]

$$
\begin{gathered}
A=\begin{array}{llllllllll}
1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\
B= & \begin{array}{lll|lll|l|l}
1 & 2 & 3 & -8 & -7 & -6 & 4 & 5
\end{array} & 9 & 10 \\
\mathrm{BP}(\mathrm{~A}, \mathrm{~B})=3
\end{array}
\end{gathered}
$$

## Quantifying Error



FN: false negative (missing edge)
$1 / 3=33.3 \%$ error rate

## NJ (BP) and NJ (INV)



120 genes, 160 leaves Uniformly Random Trees

## Additive Distance Matrix and True Evolutionary Distance (T.E.D.)



|  | S1 | S2 | S3 | S4 | S5 |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| S1 | 0 | 9 | 15 | 14 | 17 |  |
| S2 |  | 0 | 14 | 13 | 16 |  |
| S3 |  |  | 0 | 13 | 16 |  |
| S4 |  |  |  |  | 0 | 13 |
| S5 |  |  |  |  |  | 0 |

Theorem [Waterman et al. 1977] Given an m×m additive distance matrix, we can reconstruct a tree realizing the distance in $\mathrm{O}\left(\mathrm{m}^{2}\right)$ time.

## Error Tolerance of Neighbor Joining

Theorem [Atteson 1999]
Let $\left\{\mathrm{D}_{\mathrm{ij}}\right\}$ be the true evolutionary distances, and $\left\{\mathrm{d}_{\mathrm{ij}}\right\}$ be the estimated distances for T . Let $\varepsilon$ be the length of the shortest edge in $T$. If for all taxa $\mathrm{i}, \mathrm{j}$, we have

$$
\left|D_{i j}-d_{i j}\right|<\frac{1}{2} \varepsilon
$$

then neighbor joining returns T .

## BP and INV


(K: Actual number of inversions)
(Inversion-only evolution)

## Estimate True Evolutionary Distances Using BP



To use the scatter plot to estimate the actual number of events (K):

1. Compute BP/2
2. From the curve, look up the corresponding value of K
$\mathrm{BP} / 2$ vs $\mathrm{K} \quad$ (120 genes)
(K: Actual number of inversions) (Inversion-only evolution)

## Using Breakpoints to Estimate T.E.D.

- Compute $\mathrm{f}_{\mathrm{n}}(\mathrm{k})=\mathrm{E}\left[\mathrm{BP}\left(\mathrm{G}_{0}, \mathrm{G}_{\mathrm{k}}\right)\right]$ (i.e. the expected number of breakpoints after k random events; n is the number of genes)
- Given two genomes G and G':
- Compute breakpoint distance d=BP(G,G')
- Find $k$ so that $f_{n}(k)$ is closest to $d$
- Challenge: finding $f_{n}(k)$


## True Evolutionary Distance (t.e.d.) Estimators for Gene Order Data

| T.E.D. <br> Estimator | Exact-I EBP <br> [WABI'01] | Approx-I EBP <br> [STOC'01] | EDE <br> [ISMB'01] |
| :---: | :--- | :--- | :--- |
| Based on the <br> Expectation of | Breakpoint <br> distance <br> (Exact) | Breakpoint <br> distance <br> (Approx.) | Inversion <br> distance <br> (Approx.) |
| Derivation | Analytical | Analytical | Empirical |
| Model <br> knowledge | Required | Required | Inversion- <br> only |

IEBP: Inverting the Expected BreakPoint distance EDE: Empirically Derived Estimator

## Exact-I EBP [WABI'01]

- Breakpoints are identically distributed: use linearity



## State Notation

- The sign and position of gene 2 with respect to gene 1 (at pos 1 ) is $\{-n,-(n-1), \ldots,-2,2,3, \ldots, n\}$.



## Markov Chain for a Breakpoint

- Let n be the number of genes
- Each breakpoint (in particular, bp between genes 1 and 2) is a Markov process with 2(n-1) states
- We have

$$
\begin{aligned}
M_{u, v} & =(1-\alpha-\beta)\left(M_{I}\right)_{u, v}+\alpha\left(M_{T}\right)_{u, v}+\beta\left(M_{V}\right)_{u, v} \\
& =\frac{1-\alpha-\beta}{\binom{n}{2}} \iota_{n}(u, v)+\frac{\alpha}{\binom{n}{3}} \tau_{n}(u, v)+\frac{\beta}{3\binom{n}{3}} \nu_{n}(u, v)
\end{aligned}
$$

where

- $\iota_{n}(u, v)$ is the number of inversions,
- $\tau_{n}(u, v)$ is the number of transpositions,
- $\nu_{n}(u, v)$ is the number of inverted transpositions, that bring gene 2 in state $u$ to state $v$ ( $n$ is the number of genes in each genome).
- The probability trasitional matrix is easily obtained:

$$
\begin{aligned}
& \iota_{n}(u, v)= \begin{cases}\min \{|u|-1,|v|-1, n+1-|u|, n+1-|v|\} \\
0 & \text { (if } u v<0) \\
\binom{|u|-1}{2}+\binom{n+1-|u|}{2} & \text { (if } u \neq v, u v>0)\end{cases} \\
& \tau_{n}(u, v)= \begin{cases}0 & (\text { if } u=v) \\
(\min \{|u|,|v|\}-1)(n+1-\max \{|u|,|v|\}) \\
\binom{n+1-|u|}{3}+\binom{|u|-1}{3} & (\text { if } u \neq v, u v>0)\end{cases} \\
& \nu_{n}(u, v)= \begin{cases}(n-2) \iota_{n}(u, v) & \text { (if } u=v) \\
\tau_{n}(u, v) & \text { if } u \neq v=0) \\
3 \tau_{n}(u, v) & =0 v>0)\end{cases}
\end{aligned}
$$

$$
\begin{aligned}
& \begin{array}{lllllllllllllllll}
-10 & -9 & -8 & -7 & -6 & -5 & -4 & -3 & -2 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\
10
\end{array}
\end{aligned}
$$

## Exact-IEBP

- There are 2(n-1) states.
- The transitional matrix has dimension $2(n-1) \times 2(n-1)$.
- To compute $\mathrm{E}\left[\mathrm{BP}\left(\mathrm{G}_{0}, \mathrm{G}_{\mathrm{k}}\right)\right]$ for $k$ up to 2 n takes $\mathrm{O}\left(\mathrm{n}^{3}\right)$ time. (2n matrix-vector multiplications)


## Reducing the State Space



## Lower and Upper Bounds

- Under the GNT model, $s$ is constant
- $u$ is not constant, but has good lower and upper bounds: $u_{\max }$ and $u_{\text {min }}$
- Parameter $u$ is small with respect to $s$



## I nversion-Only Evolution

- Unsigned genome: $\mathrm{u}_{\min }=\mathrm{u}_{\max }->$ Markov Process [Caprara \& Lancia, 2000]
- Signed genome:

$$
\begin{gathered}
s=\frac{n-1}{\binom{n}{2}}=\frac{2}{n} \\
u_{\min }=0, u_{\max }=\frac{1}{\binom{n}{2}}
\end{gathered}
$$



- The two Markov chains ( $\mathrm{s}, \mathrm{u}_{\min }$ ) and ( $\mathrm{s}, \mathrm{u}_{\max }$ ) give lower and upper bounds to the expectation of breakpoint distance.


## GNT Model

- $s=(1-\alpha-\beta) s_{I}+\alpha s_{T}+\beta s_{I T}$

$$
\begin{aligned}
& u_{\min }=(1-\alpha-\beta) u_{I, \min }+\alpha u_{T, \min }+\beta u_{I T, \min } \\
& u_{\max }=(1-\alpha-\beta) u_{I, \max }+\alpha u_{T, \max }+\beta u_{I T, \max }
\end{aligned}
$$

- $P_{k}^{L} \leq \operatorname{Pr}\left(B_{1}\left(G_{k} \mid G_{0}\right)=1\right) \leq P_{k}^{H}$, where

$$
P_{k}^{L}=s \frac{1-\left(1-s-u_{\max }\right)^{k}}{1-\left(1-s-u_{\max }\right)} \quad P_{k}^{H}=s \frac{1-\left(1-s-u_{\min }\right)^{k}}{1-\left(1-s-u_{\min }\right)}
$$

- $\mathcal{F}_{k}=\frac{n}{2}\left(P_{k}^{L}+P_{k}^{H}\right) \sim E\left[B P\left(G_{k}, G_{0}\right)\right]$


## Approx-I EBP <br> [Wang \& Warnow, STOC'01]

Theorem Let $G_{k}$ be the genome obtained after applying $k$ random rearrangement events to genome $G_{0}$ according to the GNT model with parameters $\alpha$ and $\beta$. Let $\mathcal{F}_{k}$ be the estimate to $E\left[B P\left(G_{k}, G_{0}\right)\right]$ in the Approx-IEBP distance.
For all $k>0$,

$$
\begin{aligned}
& \left|\mathcal{F}_{k}-E\left[B P\left(G_{k}, G_{0}\right)\right]\right| \leq 1+\frac{1}{n-1}, \text { and } \\
& \phi^{-1} \leq \frac{\mathcal{F}_{k}}{E\left[B P\left(G_{k}, G_{0}\right)\right]} \leq \phi
\end{aligned}
$$

where $\phi=1+\frac{2+4 \alpha+2 \beta}{2+\alpha+\beta} n^{-1}+O\left(n^{-2}\right)$.

## True Evolutionary Distance Estimators




BP vs K
(120 genes) Exact-IEBP vs K
(K: Actual number of inversions)
(Inversion-only evolution)

## Variance of True Evolutionary Distance Estimators

- There are new distance-based phylogeny reconstruction methods (though designed for DNA sequences)
- Weighbor [Bruno et al. 2000]
uses the variance of good t.e.d.s, and yield more accurate trees than NJ .
- Variance estimates for the t.e.d.s [Wang WABI'02]

- Weighbor(IEBP), Weighbor(EDE)

K vs Exact-IEBP (120 genes)

## Deriving Var(BP)

- Difficulties in deriving $\operatorname{Var}(\mathrm{BP})$ :
- Even $E(B P)$ is only in the form of unsimplified sums [RECOMB ‘99, WABI ‘01].
- Breakpoints are not independent.
- We will use an approximating model to examine all breakpoints simultaneously
- Idea: once two adjacent genes are separated, it is hard to bring the two genes back again (especially when there are many genes).


## Approximating Model

- Approximating box model: boxes correspond to breakpoints.
- An approximation (using $n$ boxes) can be obtained in the following way:
- Every inversion chooses two boxes and put a ball in them if they are empty.
- The BP distance is approximated by the number of nonempty boxes.



## Approximating Model

- Notations:
- Let $B_{i}=1$ if box $i$ is not empty, 0 if it is.
- We use inversion-only model to illustrate; let i and j be the two breakpoints corresponding to the two endpoints of the inversion being applied.
- Let the number of breakpoints be b.
- Let n be the number of genes.


## Why the Approximation Works

- Case analysis: [Hannenhalli and Pevzner 1995]

| Case | ? BP | Condition | \# inversions |  |
| :---: | :---: | :---: | :---: | :---: |
| 1 | +2 | $\mathrm{B}_{\mathrm{i}}=\mathrm{B}_{\mathrm{j}}=0$ | $\binom{n-b}{2}$ |  |
| 2 | +1 | $\mathrm{B}_{\mathrm{i}}=0, \mathrm{~B}_{\mathrm{j}}=1$ or $\mathrm{B}_{\mathrm{i}}=1, \mathrm{~B}_{\mathrm{j}}=0$ | $b(n-b)$ |  |
| 3a | 0 | $\mathrm{B}_{\mathrm{i}}=\mathrm{B}_{\mathrm{j}}=1$ |  | Total |
| 3b | -1 -2 | $B_{i}=B_{j}=1$, one/both of $\left(g_{i-1},-g_{j}\right),\left(-g_{i}, g_{j}\right)$ adjacencies are in $\mathrm{G}_{0}$. | $\leq b$ | $\binom{b}{2}$ |

- When $b$ is small, probability of case 3 out of cases 1,2 , and 3 is small (when n is large)
- When $b$ is large, probability of $3 b / 3 c$ out of case 3 is small
- As a result we can ignore cases 3b/3c
-> As a breakpoint is asserted, it does not disappear


## Derivation of the Variance

- Fix k. Let $S=\left(\frac{1}{\binom{n}{2}}\left(x_{1} x_{2}+x_{1} x_{3}+\ldots+x_{n-1} x_{n}\right)\right)^{k}$
- Each term in the expansion of $S$ is a way of applying $k$ inversions
E. g. $x_{1}^{3} x_{2} x_{3}^{2}$ : box 1 three times, 2 once, 3 twice
- The coefficient of the term is the probabilities of such $k$ inversions
- If transpositions and inverted transpositions are present:

$$
S=\left(\frac{1-\alpha-\beta}{\binom{n}{2}} \sum_{1 \leq i<j \leq n} x_{i} x_{j}+\frac{\alpha+\beta}{\binom{n}{3}} \sum_{1 \leq i<j<l \leq n} x_{i} x_{j} x_{l}\right)^{k}
$$

- Let $S\left(a_{1}, a_{2}, \ldots, a_{n}\right)$ be the value of $S$ when we let $\mathrm{x}_{\mathrm{i}}=\mathrm{a}_{\mathrm{i}}$ for all i.
- Let $S_{j}=S(\underbrace{1,1,1, \ldots, 1}_{j 1^{\prime} s}, 0, \ldots, 0)$


## Derivation of $\operatorname{Var}(\mathbf{B P})$

- Let $u_{i}$ be the sum of coefficients of all terms in the expansion of S in the tollowing form:

$$
x_{1}^{a_{1}} x_{2}^{a_{2}} \cdots x_{i}^{a_{i}}\left(a_{1}, a_{2}, \ldots, a_{i}>0\right)
$$

Then $\binom{n}{i} u_{i}$ is the probability of having i nonempty boxes after k events.

- We want to compute

$$
Z_{a}=\sum_{i=0}^{n} i(i-1) \cdots(i-a+1)\binom{n}{i} u_{i}=n(n-1) \cdots(n-a+1) \sum_{i=a}^{n}\binom{n-a}{i-a} u_{i}
$$

In particular,

$$
\begin{aligned}
& z_{1}=\sum_{i=1}^{n} i\binom{n}{i} u_{i}=E[b \mid k] \approx E\left[B P\left(G_{0}, G_{k}\right)\right] \\
& z_{2}=\sum_{i=1}^{n} i(i-1)\binom{n}{i} u_{i}=E\left[b^{2}-b \mid k\right] \approx E\left[B P^{2}\left(G_{0}, G_{k}\right)-B P\left(G_{0}, G_{k}\right)\right]
\end{aligned}
$$

$$
\begin{aligned}
S & =\left(\frac{1}{\binom{n}{2}}\left(\sum_{1 \leq i<j \leq n} x_{i} x_{j}\right)\right)^{k} \\
& =\sum_{1 \leq i \leq n} \sum_{\left\{t_{1}, t_{2}, \ldots, t_{i}\right\} \subseteq\{1,2, \ldots, n\}} \sum_{\substack{a_{1}, a_{2}, \ldots, a_{i} \geq 1 \\
a_{1}+a_{2}+\ldots+a_{i}=2 k}} c\left(t_{1}, t_{2}, \ldots, t_{i}, a_{1}, a_{2}, \ldots, a_{i}\right) x_{t_{1}}^{a_{1}} x_{t_{2}}^{a_{2}} \cdots x_{t_{i}}^{a_{i}}
\end{aligned}
$$

$$
S_{j}=\sum_{1 \leq i \leq j} \sum_{\left\{t_{1}, t_{2}, \ldots, t_{i}\right\} \subseteq\{1,2, \ldots, j\}} \sum_{\substack{a_{1}, a_{2}, \ldots, a_{i} \geq 1 \\ a_{1}+a_{2}+\ldots+a_{i}=2 k}} c\left(t_{1}, t_{2}, \ldots, t_{i}, a_{1}, a_{2}, \ldots, a_{i}\right)
$$

$$
=\sum_{1 \leq i \leq j} \sum_{\left\{t_{1}, t_{2}, \ldots, t_{i}\right\} \subseteq\{1,2, \ldots, j\}} u_{i}=\sum_{1 \leq i \leq j}\binom{j}{i} u_{i}
$$

Lemma Let $a$ be some given integer such that $1 \leq a \leq n$. Let us be given $\left\{u_{1}, u_{2}, \ldots, u_{n}\right\}$ such that

$$
\sum_{i=0}^{j}\binom{j}{i} u_{i}=\sum_{i=0}^{n}\binom{j}{i} u_{i}=S_{j}
$$

for all $j, 1 \leq j \leq n$. We have

$$
\sum_{i=n-a}^{n}(-1)^{n-i}\binom{a}{n-i} S_{i}=\sum_{i=0}^{n}\binom{n-a}{i-a} u_{i}
$$

## Expectation and Variance [WABI'02]

- Let $b_{k}$ be the number of nonempty boxes after $k$ (box choosing) iterations in the approximation model. Let $a$ $+\beta=$ ?. We have

$$
\begin{aligned}
& S_{n-1}=\left(1-\frac{2+\gamma}{n}\right)^{k}, S_{n-2}=\left(\frac{(n-3)(n-2-2 \gamma)}{n(n-1)}\right)^{k} . \\
& E b_{k}=n\left(1-S_{n-1}\right) \\
& \operatorname{Varb}_{k}=n S_{n-1}-n^{2} S_{n-1}^{2}+n(n-1) S_{n-2}^{2}
\end{aligned}
$$

- We use the delta method to obtain the variance of IEBP:

$$
\operatorname{Var} \widehat{k}\left(b_{k}\right) \simeq\left(\frac{d}{d k} E b_{k}\right)^{-2} \operatorname{Var} b_{k}=\frac{\left(1-n S_{n-1}+(n-1)\left(\frac{S_{n-2}}{S_{n-1}}\right)\right)}{n S_{n-1}\left(\ln \left(1-\frac{2+\gamma}{n}\right)\right)^{2}}
$$

## Simulation Results


$\operatorname{Var}\left(B P_{k}\right)$
Variance of BP distance after $k$ events


$$
\operatorname{Var} \widehat{k}\left(b_{k}\right)
$$

Variance of IEBP

## Regression Formula for E(INV) and Var(INV)

- Let n be the number of genes, x be the normalized number of inversions ( $k / n$ ), and $f(x)$ be the normalized expectation of the inversion distance ( $f(x)$ seems to be roughly independent of $n$ )
- We use nonlinear regression to obtain easily computable formulas for $E(I N V)$ and $\operatorname{Var}(I N V)$ :

$$
\begin{aligned}
& f(x)=\min \left\{\frac{x^{2}+b x}{x^{2}+c x+b}, x\right\} \quad\left(x=\frac{k}{n}\right) \\
& \text { 1. } f(0)=0 \quad f^{\prime}(0)=1 \\
& \text { 3. } 0 \leq f(x) \leq x \\
& \text { 4. } f^{-1}(y) \quad \text { exists for all } y: 0 \leq y \leq 1
\end{aligned}
$$

$->\quad b=0.5956, c=0.4577$

## EDE

## [Moret, Wang, Warnow, \& Wyman, ISMB'01]



## Formula for Var(INV) and Var(EDE)

- Let n be the number of genes, x be the normalized number of inversions ( $k / n$ ), and $g_{n}(x)$ be the standard deviation of the inversion distance.
- The regression of $g_{n}(x)$ : we use the following form

$$
\begin{gathered}
g_{n}(x)=n^{q} \frac{u x^{2}+v x}{x^{2}+w x+t} \\
\mathrm{q}=-0.6998, \mathrm{u}=0.1684, \mathrm{v}=0.1573, \mathrm{w}=-1.3893 \text {, and } \\
\mathrm{t}=0.8224 .
\end{gathered}
$$

- $\operatorname{Var}(E D E)$ can be obtained using the delta method on $\operatorname{Var}(I N V)$.


## Regression for Var(INV)

Regression: solid lines, Simulation: dots


## Distance-Based Methods



## Using T.E.D. Helps



## IEBP is Robust to Model Violations



120 genes, 160 taxa Uniformly Random Trees (alpha, beta) $=(0,0)$ (inversion only)

## Maximum Parsimony Returns Thousands of Trees

- Example:
- The complete Caesalpinia dataset: 7095 trees on 82 taxa.
- The Astericeae dataset: 34,560 trees on 288 taxa.
- Consensus methods are necessary so we can summarize so many trees.
- Current approaches are limited to the strict consensus and majority consensus trees, and lose information


## Postprocessing: Traditional Approaches

- Single-tree consensus

Example: strict consensus

$\left(t_{1}, t_{2}, t_{3}\right.$ all refine $\left.t\right)$


## How Do We Interpret the Consensus Tree

- Given a nonbinary consensus tree $t$, every binary tree that refines $t$ is equally probable to be the true tree:

(15 refinement trees)


## Disadvantages of Single-Tree Consensus

- Loses a lot of information
- Sensitive to outlier trees
- Sensitive to small perturbations in the dataset


## Sometimes A Cluster is Enough (Campanulaceae)



## The Campanulaceae Gene-Order Dataset

1. 13 taxa (outgroup Tobacco)
2. 216 trees
(Courtesy Nina Amenta and Jeff Klingner)

## Complex Structure in the I nferred Set of Trees



## The Caesalpinia cpDNA Dataset

1. 51 taxa
2. 342 trees
(Courtesy Nina Amenta and Jeff Klingner)

## Why We Want to Cluster Trees

- Dividing trees into clusters, and use the consensus trees from each cluster to represent "conflicting hypotheses" for the true phylogeny.
- Merits:
- Represent the input set of trees better
- Identify outliers
- Restrict perturbations to a small number of clusters


## Biological Criteria

- Number of clusters
- Number of edges of the consensus
- Diameter of a cluster
- Density of clusters
- Etc.


## Information Loss: How We I nterpret the Clustering

- We can define distributions for both the original set of trees and the clustering.

Input set of tree $T$ :
All trees are equally probable.


Clustering_ $\left\{C_{1}, C_{2}, \ldots, C_{k}\right\}$ :
All trees refining any of $S C\left(C_{j}\right)$ are equally probable.


## Distributions

- Input set of tree T:

$$
f_{T}(t)=\left\{\begin{array}{cl}
\frac{1}{|T|} & \text { if } t \in T \\
0 & \text { othewise }
\end{array}\right.
$$

- Clustering $\left\{\mathrm{C}_{1}, \mathrm{C}_{2}, \ldots, \mathrm{C}_{\mathrm{k}}\right\}$ : let

$$
B=\bigcup_{i=1}^{k} B\left(C_{i}\right)
$$

$$
f_{C}(t)= \begin{cases}\frac{1}{|B|} & \text { if } t \in B \\ 0 & \text { otherwise }\end{cases}
$$

(Here $B(C)$ is the set of binary trees that refine the strict consensus of C )

## Information Loss (KL)

- The distance between the two distributions is the loss of information due to clustering.
- $\mathrm{L}_{1}$ distance
- $\left.\begin{array}{ll}\mathrm{L}_{2} & \text { distance } \\ \text { - } \mathrm{L}_{\infty} \text { distance }\end{array}\right\} L_{x}(T, C)=\sum_{t}\left\|f_{T}(t)-f_{C}(t)\right\|_{x}$
- Kullback-Leibler distance (relative entropy):

$$
K L(T, C)=\sum_{t} f_{T}(t) \ln \frac{f_{T}(t)}{f_{C}(t)}
$$

## Postprocessing of Phylogenetic Analysis Using Clustering [ISMB'02]

- The first framework using clustering algorithms in the postprocessing of phylogenetic analyses.
- Improves upon the traditional single-consensus approach in terms of information loss
- Identifies outliers in the Caesalpinia dataset
- Improves the resolution of the strict consensus by $36 \%$
- Only loses 4\% of the trees


Number of Clusters

## Caesalpinia (51 taxa, 450 trees)

| Clu No. | No. of Trees | \% Edges lost |
| :---: | :---: | :---: |
| lclu | 450 | $22.9 \%$ |
| 1 | 108 | $10.4 \%$ |
| 2 | 324 | $12.5 \%$ |
| 3 | 18 | $10.4 \%$ |
| $1+2$ | 432 | $14.6 \%$ |

$\mathrm{KL}($ Agg-complete, 3 clu$)=1.449269$
$\mathrm{KL}(1 \mathrm{clu})=9.790346$
Improvement: (22.9-14.6)/22.9 $=36 \%$
$\%$ trees dropped: $18 / 450=4 \%$

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## - Genome rearrangement phylogeny

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## Postprocessing by clustering

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