

# Algorithms and geometric constructions

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- 1930s: Gödel, Kleene, Church, Turing, Post defined representative classes of algorithms thus making possible the proofs of algorithmic undecidability

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- our topic: the evolution of a notion of geometric construction (special class of algorithms)

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- but what is the problem?

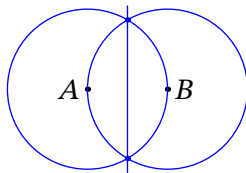
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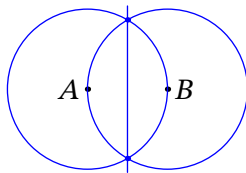
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- operations:
  - (ruler) draw a line through two points;
  - given three points  $A, B, C$ , draw a circle with center  $A$  and radius  $BC$ ;
  - add the intersection points of existing curves

Define a configuration to be a finite collection  $\mathcal{C}$  of points, lines, and circles in the Euclidean plane. Define a construction step to be one of the following operations to enlarge the collection  $\mathcal{C}$ :

- (Straightedge) Given two distinct points  $A, B$  in  $\mathcal{C}$ , form the line  $\overline{AB}$  that connects  $A$  and  $B$ , and add it to  $\mathcal{C}$ .
- (Compass) Given two distinct points  $A, B$  in  $\mathcal{C}$ , and given a third point  $O$  in  $\mathcal{C}$  (which may or may not equal  $A$  or  $B$ ), form the circle with centre  $O$  and radius equal to the length  $|AB|$  of the line segment joining  $A$  and  $B$ , and add it to  $\mathcal{C}$ .
- (Intersection) Given two distinct curves  $\gamma, \gamma'$  in  $\mathcal{C}$  (thus  $\gamma$  is either a line or a circle in  $\mathcal{C}$ , and similarly for  $\gamma'$ ), select a point  $P$  that is common to both  $\gamma$  and  $\gamma'$  (there are at most two such points), and add it to  $\mathcal{C}$ .

We say that a point, line, or circle is constructible by straightedge and compass from a configuration  $\mathcal{C}$  if it can be obtained from  $\mathcal{C}$  after applying a finite number of construction steps.

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- “points for which we do not have affine or metric information” (Bieberbach)
- “What we do not include in our analysis are arbitrary elements that are used in some constructions. Their role is not so simple, however, as it is sometimes thought” (Tietze)

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- definition: *a problem (for a given input configuration) is solvable by ruler and compass = Alice has a winning strategy in the corresponding game*

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  - $\mathcal{L}$  contains all the input objects;
  - $\mathcal{L}$  is closed under three basic operations (line through given points, circle with given center and radius, intersection points)
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- another direction: essentially proven by Akopyan and Fedorov (though they avoid using formal definitions!)



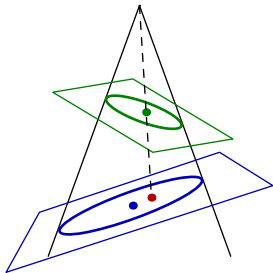
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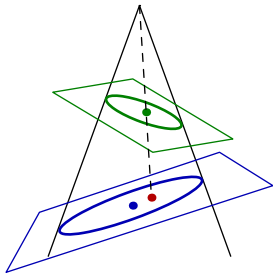
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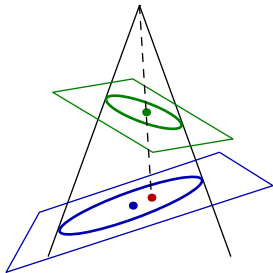
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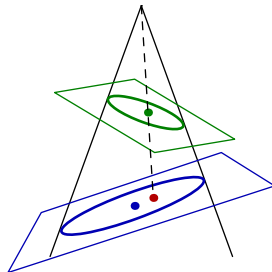
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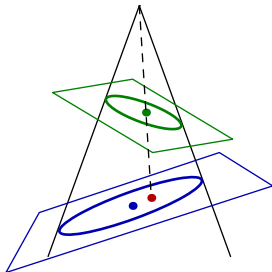
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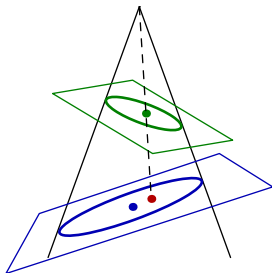
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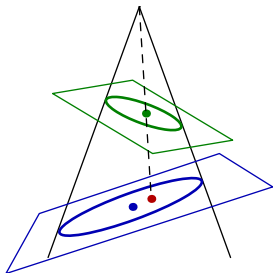
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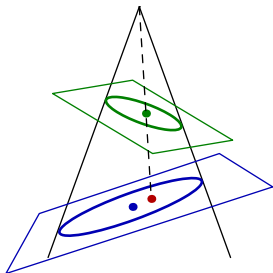
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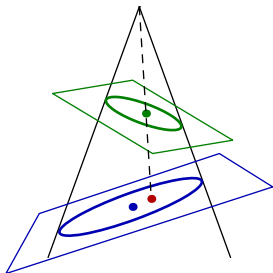
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- proof does not work, but the statement is true (for the game definition); it is important that many projective transformations exist
- Cauer's "theorem" is plainly false (Akopyan–Fedorov): for some pairs of disjoint circles one *can* construct the center (but for other pairs one cannot)



# Conclusions

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