Crypto 2022. Final exam preparation.

Shannon's entropy

Exercise 1. Prove that for every triple of jointly distributed random variables (X, Y, Z) we have

(a) $H(X, Y|Z) \le H(X|Z) + H(Y|Z)$, (b) $2H(X, Y, Z) \le H(X, Y) + H(X, Z) + H(Y, Z|X)$, (c) $H(X) \le H(X|Y) + H(X|Z) + I(Y : Z)$, (d) $H(X|Z) \le H(X|Y) + H(Y|Z)$.

Exercise 2. There is a pair of random variables X and Y that are distributed on the set $\{1, 2, ..., n\}$. It is known that $\operatorname{Prob}[X \neq y] = \epsilon$. Prove that

$$H(X|Y) \le 1 + \epsilon \log(n-1).$$

Hint: It is useful to introduce a new variables

$$Z := \begin{cases} 1, & \text{if } X = Y, \\ 0, & \text{otherwise} \end{cases}$$

and compare H(X|Y) with H(X|Y,Z) + H(Z).

Secret sharing

Exercise 3. We need to share a secret k (which is a bit string of length n) among four participants Alice, Bob, Charlie, Dan in such a way that the minimal groups that known the secret are {Alice, Bob}, {Alice, Charlie}, {Alice, Dan} (Alice alone as well as Bob, Charlie and Dan together should have no information on the secret).

(a) Construct a secret sharing scheme with the required property.

(b) Construct a secret sharing scheme with the required property so that each share of the secret S_A , S_B , S_C , S_D is a represented by a string of n bits.

(c) Show that in every secret sharing scheme for Alice, Bob, Charlie, Dan the space of shares for Alice consists of at least 2^n different values (i.e., we cannot give to Alice less than n bits of information.

Exercise 4. We need to share a secret k (which is a bit string of length n) among four participants Alice, Bob, Charlie, Dan in such a way that the minimal groups that known the secret are

{Alice, Bob}, {Alice, Charlie}, {Alice, Dan}, {Bob, Charlie, Dan}.

(a) Construct a secret sharing scheme with the required property.

(b)* Show that in this case in every secret sharing scheme one of the participants must receive a share with strictly more than 2^n possible values (i.e., we cannot give to each participant only n bits of information),

Shamir's secret sharing and polynomials the modular arithmetic.

Exercise 5. Let a_0 , a_1 , a_2 , a_3 be integer numbers, and

$$f(x) = a_0 + a_1 x + a_2 x^2 + a_3 x^3.$$

It is know that

 $f(1) = 0 \mod 31, f(2) = 0 \mod 31, f(3) = 0 \mod 31, f(4) = 30 \mod 31.$

Find $a_0 \mod 31$.

Exercise 6. Find a triple of numbers a_0, a_1, a_2 (not all congruent to zero modulo 10) such that the polynomial $f(x) = a_0 + a_1x + a_2x^2$ has at least three different roots in $\mathbb{Z}/10\mathbb{Z}$.

Exercise 7. For the following n find without a computer a number k such that $100 \cdot k = 1 \mod n$.

(a) n = 61; (b) n = 599; (c) n = 1009.

Exercise 8. It is known that the number $p = 5\,104\,051$ is prime.

(a) Show without using a computer that the number $g_1 = 3$ the numbers

$$g, (g^2 \mod p), (g^3 \mod p), \ldots$$

cover the whole set $\{1, 2, \ldots, p-1\}$. *Hint:* compute $g_1^{(p-1)/2} \mod p$.

(b) Show without using a computer that the same property is not true for the number $g_2 = 9$.

(c) Let $g_3 = 522904$. Count the number of different elements in the list

$$\{1, 2, \ldots, p-1\}$$

(d) Using (c), find without a computer a number x such that

$$g_3^x = 4581146 \mod p.$$

Moralité de la fable : if the subgroup generated by g is small, then the problem of discrete logarithm with the base g is simple.