Single-exponential algorithms and linear kernels via protrusion decompositions

Eun Jung Kim[1] Christophe Paul[2] Ignasi Sau[2]

Alexander Langer[3] Felix Reidl[3] Peter Rossmanith[3] Somnath Sikdar[3]

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[1] CNRS, LAMSADE, Paris (France)

[2] CNRS, LIRMM, Montpellier (France)

[3] Department of Computer Science, RWTH Aachen University (Germany)

Outline of the talk

1. Preliminaries
2. Protrusion decompositions
   - Definitions
   - A simple algorithm to compute them
3. Single-exponential algorithm for Planar-$\mathcal{F}$-Deletion
   - Motivation and our result
   - Sketch of proof
   - Further research
4. Linear kernels on graphs without topological minors
   - Motivation and our result
   - Idea of proof
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Some words on parameterized complexity

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$$f(k) \cdot n^{O(1)}, \quad \text{for some function } f.$$

  Examples: k-Vertex Cover, k-Longest Path.

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- A single-exponential parameterized algorithm is an FPT algo s.t.

$$f(k) = 2^{O(k)}.$$

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- Linear-time algorithms based on modular decompositions.

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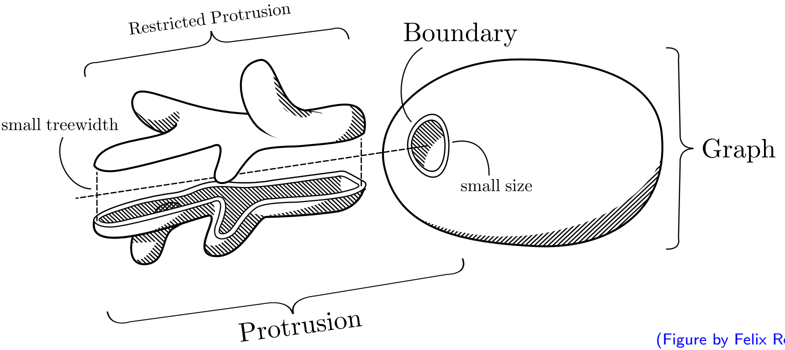
Protrusions

[Bodlaender, Fomin, Lokshtanov, Penninkx, Saurabh, Thilikos '09]

- Given a graph G, a set $W \subseteq V(G)$ is a t-protrusion of G if

$$|\partial_G(W)| \leqslant t \quad \text{and} \quad \mathrm{tw}(G[W]) \leqslant t$$



(Figure by Felix Reidl)

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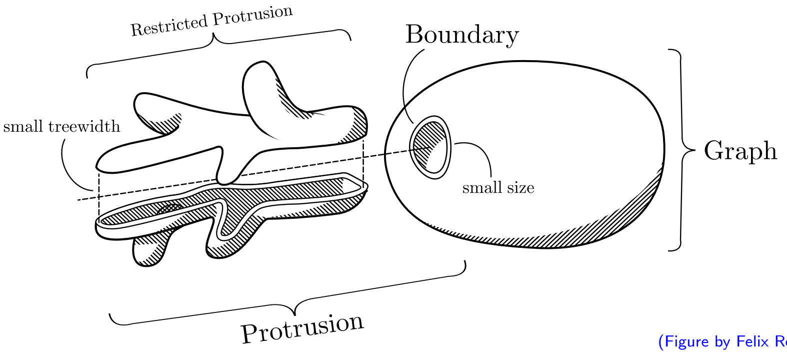


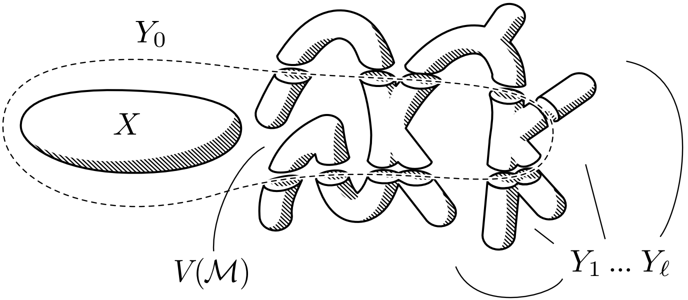
(Figure by Felix Reidl)

- The vertex set $W' = W \setminus \partial_G(W)$ is the restricted protrusion of W.
- We call $\partial_G(W)$ the boundary and $|W|$ the size of W.

Protrusion decompositions

An (α, t)-protrusion decomposition of a graph G is a partition $\mathcal{P} = Y_0 \uplus Y_1 \uplus \cdots \uplus Y_\ell$ of $V(G)$ such that:

- for every $1 \leqslant i \leqslant \ell$, $N(Y_i) \subseteq Y_0$;
- for every $1 \leqslant i \leqslant \ell$, $Y_i \cup N_{Y_0}(Y_i)$ is a t-protrusion of G;
- $\max\{\ell, |Y_0|\} \leqslant \alpha$.



The set Y_0 is called the separating part of $\mathcal{P}$.

(Figure by Felix Reidl)

Next subsection is...

Main (informal) ideas of our algorithm

- Protrusion decompositions have already been used in the literature.

[Bodlaender, Fomin, Lokshtanov, Saurabh, Thilikos '09-12]

Main (informal) ideas of our algorithm

- Here we present a new algorithm to compute protrusion decompositions for graphs G that come equipped with a set

  $$\boxed{X \subseteq V(G) \text{ s.t. } \mathrm{tw}(G - X) \leqslant t}$$

  for some constant $t > 0$.

  The set X is called a t-treewidth-modulator.

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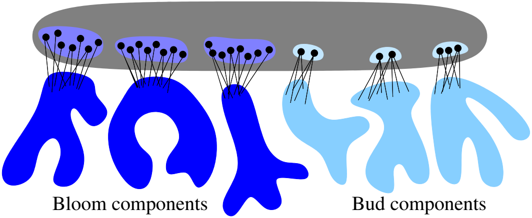
- Given tree-decompositions of the conn. comp. of $G - X$ with $\geqslant r$ neighbors in X, we identify a set of bags $\mathcal{M}$ in a bottom-up manner.



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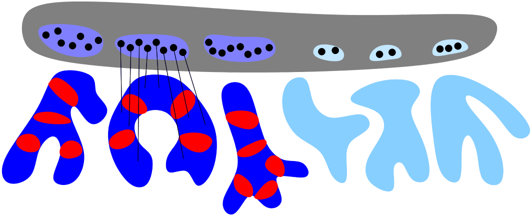

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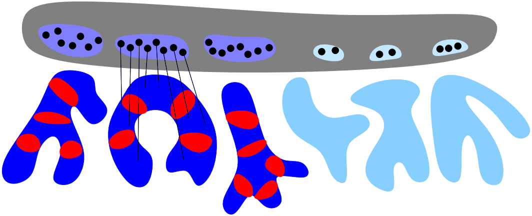

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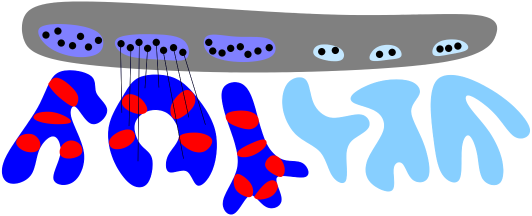

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- Some marked bags will be mapped bijectively into pairwise vertex-disjoint connected subgraphs of $G - X$, each of which has $\geqslant r$ neighbors in X.
- Finally, to guarantee that the conn. comp. of $G - (X \cup V(\mathcal{M}))$ form protrusions with small boundary, the set $\mathcal{M}$ is closed under taking LCA.

Description of the bag marking algorithm

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    if B is the LCA of two marked bags of $\mathcal{M}$:

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Return $Y_0 = X \cup V(\mathcal{M})$.

Some properties of the bag marking algorithm

Lemma

The bag marking algorithm can be implemented to run in $O(n)$ time, where the hidden constant depends only on t and r.

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Given a graph G and a subset $S \subseteq V(G)$, a cluster of $G - S$ is a maximal collection of connected components of $G - S$ with the same neighborhood in S.

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Proposition

- *Let* r, t *be two positive integers,*
- *let* G *be a graph and* $X \subseteq V(G)$ *such that* $\mathrm{tw}(G - X) \leqslant t$*,*
- *let* $Y_0 \subseteq V(G)$ *be the output of the algorithm with input* (G, X, r)*, and*
- *let* $Y_1, \ldots, Y_\ell$ *be the set of clusters of* $G - Y_0$*.*

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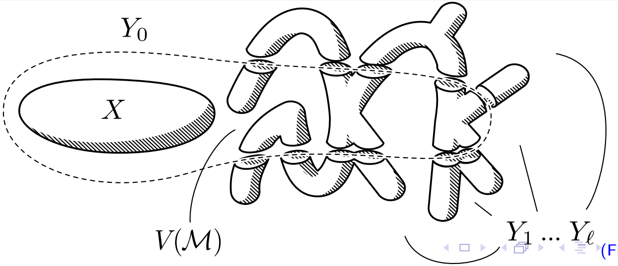
- *let $Y_1, \ldots, Y_\ell$ be the set of clusters of $G - Y_0$.*

Then $\mathcal{P} := Y_0 \uplus Y_1 \uplus \cdots \uplus Y_\ell$ is a $(\max\{\ell, |Y_0|\}, 2t + r)$-protrusion decomp. of G.



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Some particular cases:

1. $\mathcal{F} = \{K_2\}$: $\equiv$ VERTEX COVER
   $\equiv$ TREEWIDTH-ZERO VERTEX DELETION
2. $\mathcal{F} = \{K_3\}$: $\equiv$ FEEDBACK VERTEX SET
   $\equiv$ TREEWIDTH-ONE VERTEX DELETION
3. $\mathcal{F} = \{K_4\}$: $\equiv$ TREEWIDTH-TWO VERTEX DELETION

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- $\mathcal{F} = \{K_2\}$ $O^*(1.2738^k)$ [Chen, Fernau, Kanj, Xia '10]
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- $\mathcal{F} = \{\theta_r\}$ $O^*(c^k)$ [Joret, Paul, S., Saurabh, Thomassé '11]
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- $2^{O(k \log k)} \cdot n^2$ -time algorithm. [Fomin, Lokshtanov, Misra, Saurabh '11]

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- $2^{O(k \log k)} \cdot n^2$ -time algorithm. [Fomin, Lokshtanov, Misra, Saurabh '11]
- $2^{O(k)} \cdot n \log^2 n$ -time algorithm for PLANAR-CONNECTED-$\mathcal{F}$-DELETION. [Fomin, Lokshtanov, Misra, Saurabh '12]

Our result

Theorem

The PLANAR-$\mathcal{F}$-DELETION *problem can be solved in time* $2^{O(k)} \cdot n^2$.

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- This result unifies a number of algorithms in the literature.
- No hope for a $2^{o(k)} \cdot n^{O(1)}$-time algorithm (under ETH). [Chen *et al.* '05]

  That is, the function $2^{O(k)}$ in our theorem is best possible.

Next subsection is...

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Input: A graph G, a non-negative integer k, and a set $X \subseteq V(G)$ with $|X| = k$ s.t. $G - X$ is $\mathcal{F}$-minor-free.

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Lemma (well-kwown)

If DISJOINT PLANAR-$\mathcal{F}$-DELETION *can be solved in time* $O^*(c^k)$ *for some* $c \in \mathbb{N}^+$*, then* PLANAR-$\mathcal{F}$-DELETION *can be solved in* $O^*((c+1)^k)$.

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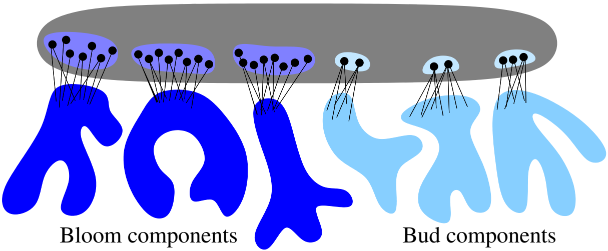
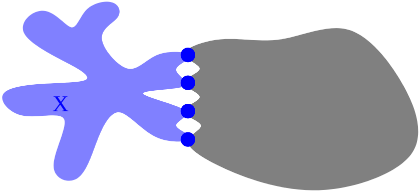


Linear protrusion decompositions

⋆ Recall that a β-protrusion in a graph G is a subset $Y \subseteq V(G)$ such that

$$|\partial(Y)| \leqslant \beta \text{ and } \mathrm{tw}(G[Y]) \leqslant \beta$$



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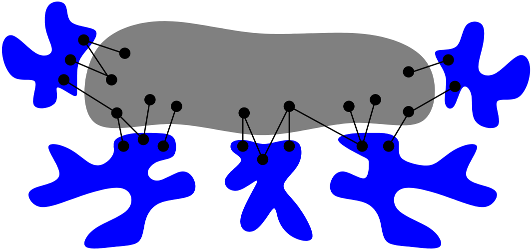

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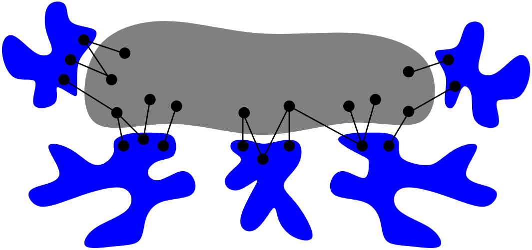

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⋆ $\mathcal{P}$ is linear with respect to a parameter k whenever $\alpha = O(k)$.

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- We will use our algorithm to compute protrusion decompositions.

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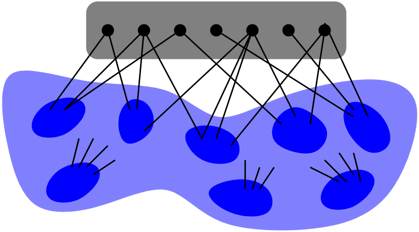
⋆ Both steps can be done in single-exponential time.

First step: analysis of the bag marking algorithm

Lemma (edge simulation to chop bloom components)

If $C_1, \ldots, C_\ell$ is a collection of connected pairwise vertex-disjoint subgraphs of $G - X$ such that $|N_X(C_i)| \geqslant r$ for $1 \leqslant i \leqslant \ell$, then $\boxed{\ell \leqslant (1 + \alpha_r) \cdot k}$.



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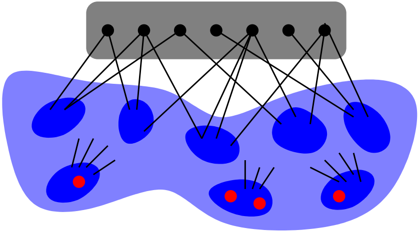

Proposition (Thomason '01)

There exists a constant $\alpha < 0.320$ such that any n-vertex graph with no K_r-minor has at most $\alpha_r \cdot n = (\alpha \cdot r\sqrt{\log r}) \cdot n$ edges.

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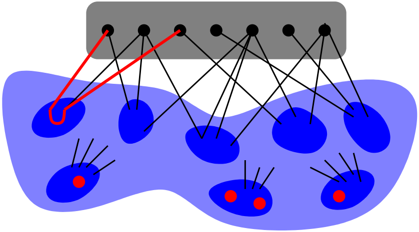

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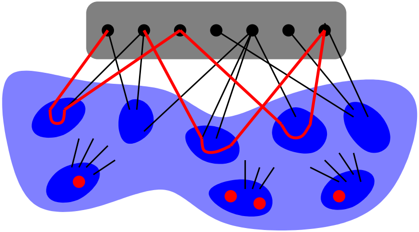

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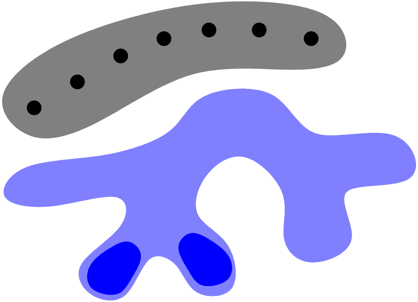

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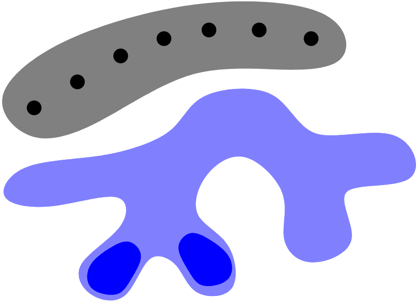

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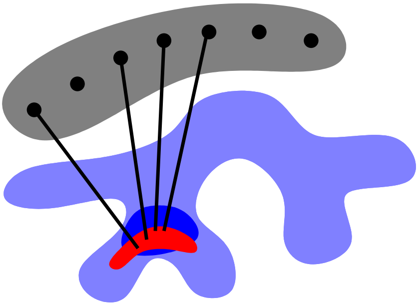

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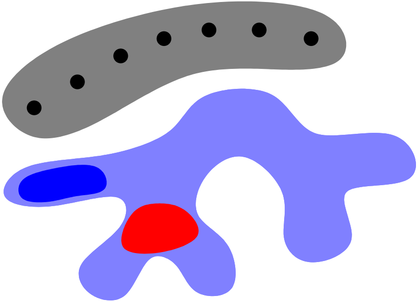

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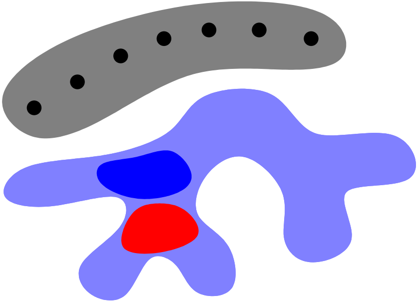

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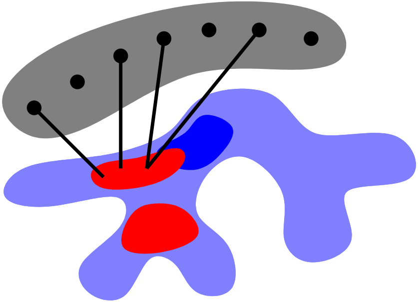

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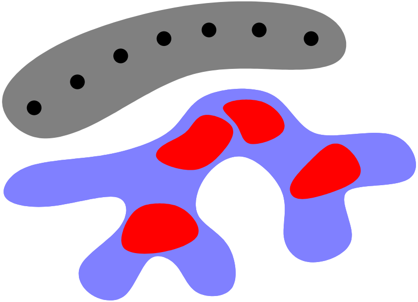

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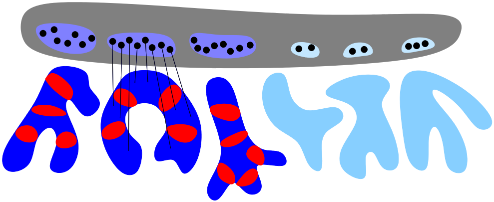

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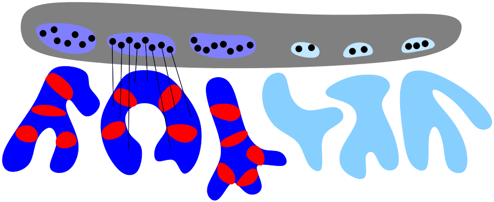

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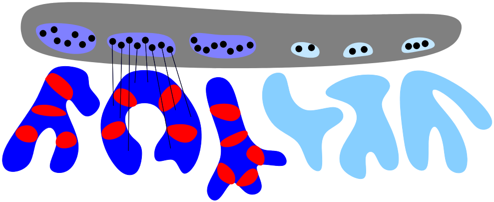

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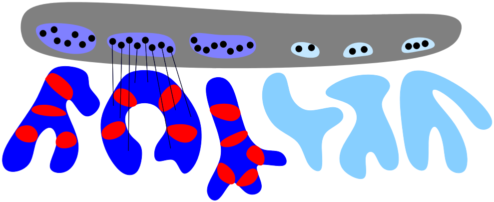

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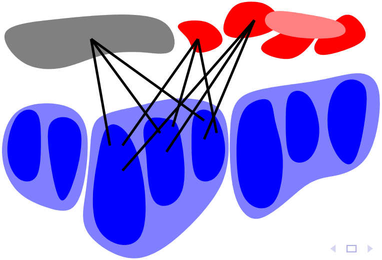

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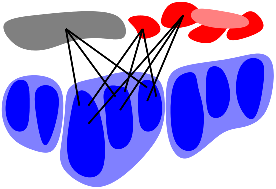

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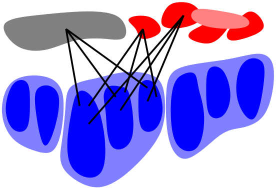

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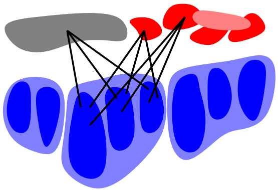

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⋆ At most $\ell' = k - |I|$ clusters $C_1, \ldots, C_{\ell'}$ intersect the alternative solution $\tilde{X}$.

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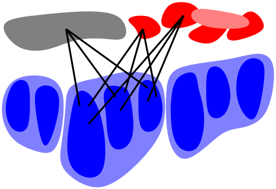

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We have that $G' = G_I - \cup_{i=1}^{\ell'} C_i$ is $\mathcal{F}$-minor-free.

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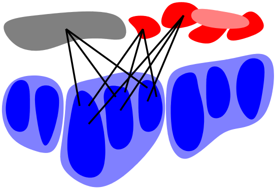

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⋆ Using edge simulation we construct a minor of G' on vertices of Y_0.

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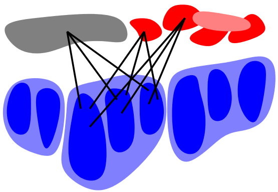

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⋆ As before, the number of clusters used so far is at most $\alpha_r \cdot k$.

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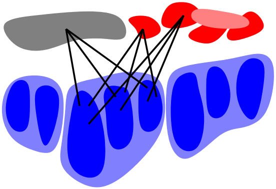

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⋆ When we cannot add more edges, all neighborhoods of clusters are cliques!

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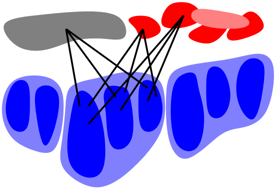

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⋆ Now we use the Proposition: the number of remaining clusters is $\mu_r \cdot k$.

Back to the road map of the algorithm

Therefore, the partition $\mathcal{P} = Y_0 \uplus C_1 \uplus \cdots \uplus C_\ell$ is a

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Recall the two main steps of our algorithm:

1 Guess the intersection $I = \tilde{X} \cap Y_0$ of the alt. solution $\tilde{X}$ with Y_0 s.t.:

- $G - I$ has a linear protrusion decomposition

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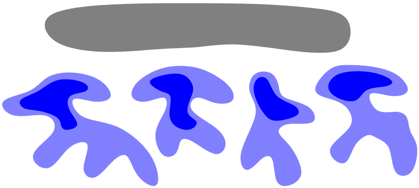
2 Finally, compute $\tilde{X} \setminus I$, given a linear protrusion decomposition.

Based on the finite index of MSO-definable properties (automaton theory)

Solving the problem when given a linear protrusion decomposition



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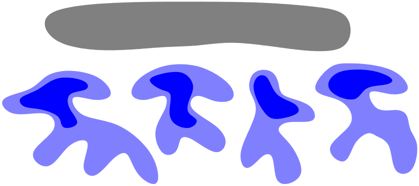
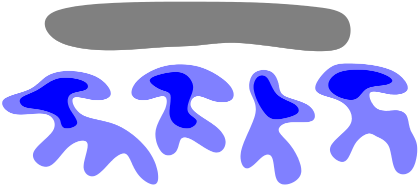


Main ingredients of our approach:

- ⋆ We define an equivalence relation on subsets of vertices of each restricted protrusion Y_i (roughly, same class if they behave in the same way).



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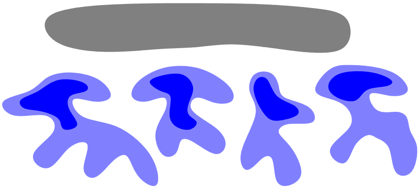

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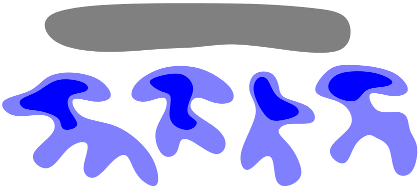

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- ⋆ We use a decomposability property of the solution: there exists a solution which is formed by the union of one representative per restricted protrusion.
- ⋆ To make the algorithm constructive and uniform on the family $\mathcal{F}$, we use classic arguments from tree automaton theory (like method of test sets).

Next subsection is...

Conclusions and further research

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For $\mathcal{F} = \{K_5, K_{3,3}\}$, an explicit FPT algorithm is known.
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★ We could forbid the family of graphs $\mathcal{F}$ according to another containment relation, like topological minor.

Next section is...

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Kernels

- A kernel for a parameterized problem Π is an algorithm that given (x, k) outputs, in time polynomial in $|x| + k$, an instance (x', k') s.t.:
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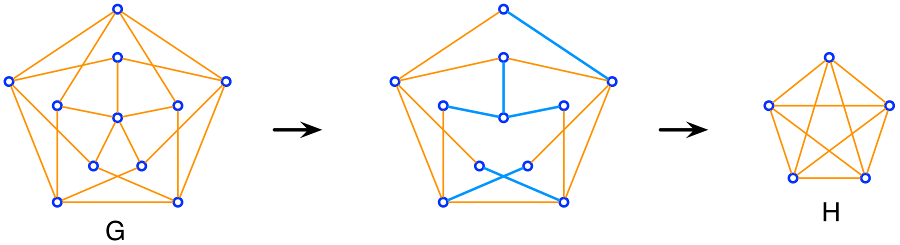
- Question: which FPT problems admit linear or polynomial kernels?

Minors and topological minors



- *H* is a minor of a graph *G* if *H* can be obtained from a subgraph of *G* by contracting edges.

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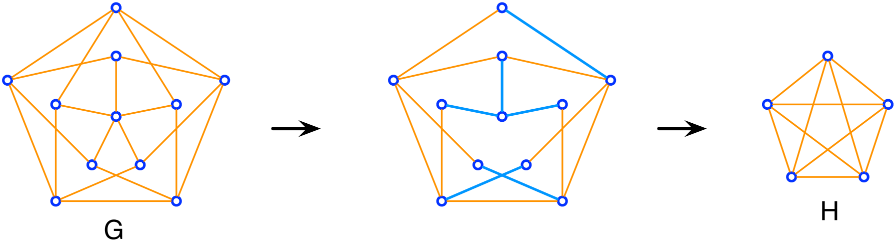

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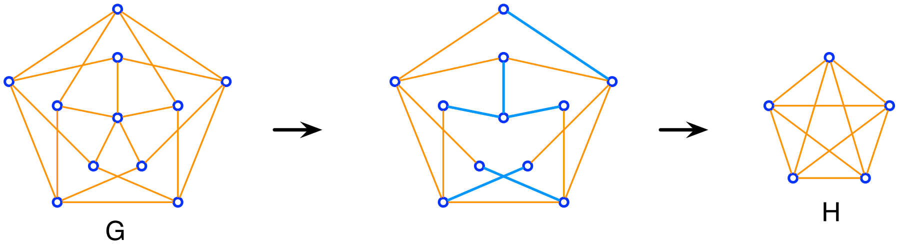

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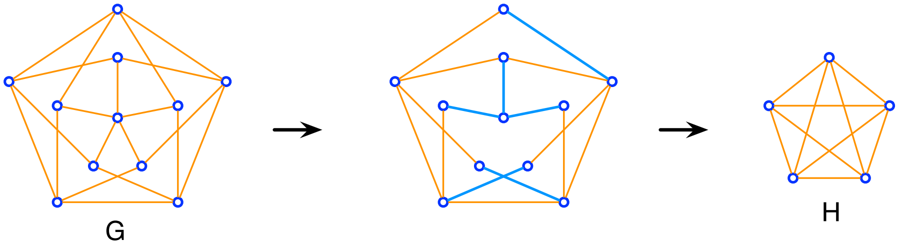

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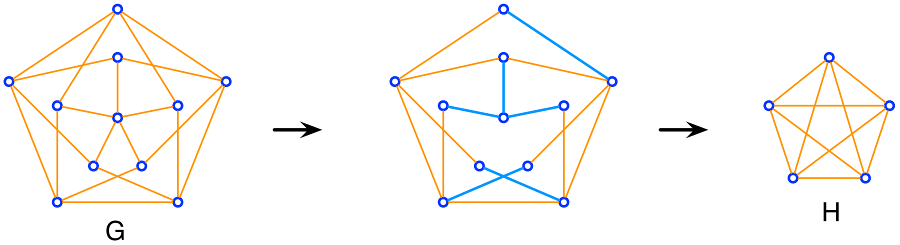

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- Fixed H: H-minor-free graphs $\subseteq$ H-topological-minor-free graphs.

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- Meta-result for H-topological-minor-free graphs. [Our result]

Our result

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Fix a graph H. Let Π be a parameterized graph problem on the class of H-topological-minor-free graphs that is treewidth-bounding and has finite integer index (FII). Then Π admits a linear kernel.

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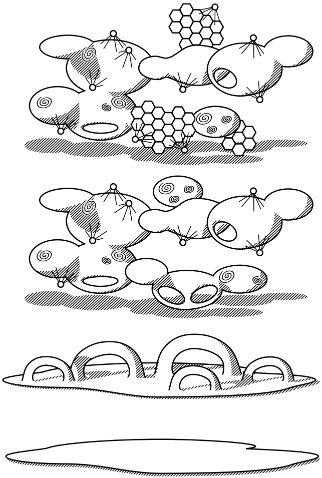
Problems affected by our result:

Treewidth-t Vertex Deletion, Chordal Vertex Deletion, Interval Vertex Deletion, Edge Dominating Set, Feedback Vertex Set, Connected Vertex Cover, ...

Linear kernels on sparse graphs – the conditions



H-topological-minor-free	treewidth-bounding
$\bigcup$	
H-minor-free	bidimensional, separation property
$\bigcup$	
bounded genus	quasi-compact
$\bigcup$	
planar	“distance-property”

(Figure by Felix Reidl)

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- Thus, our results imply the linear kernels of [Fomin, Lokshtanov, Saurabh, Thilikos '10]

Next subsection is...

Finite Integer Index (FII)

[Bodlaender, de Fluiter '01]

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- Let Π be a parameterized graph problem restricted to a graph class $\mathcal{G}$ and let G_1, G_2 be two *t*-boundaried graphs in $\mathcal{G}_t$.

Finite Integer Index (FII)

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- Let Π be a parameterized graph problem restricted to a graph class $\mathcal{G}$ and let G_1, G_2 be two t-boundaried graphs in $\mathcal{G}_t$.
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- The protrusion limit of Π is a function $\rho_\Pi \colon \mathbb{N} \to \mathbb{N}$ defined as $\rho_\Pi(t) = \max_{G \in \mathcal{R}_t} |V(G)|$. We also define $\rho'_\Pi(t) = \rho_\Pi(2t)$.

Disconnected Planar-$\mathcal{F}$-Deletion has not FII

- We prove: if $\mathcal{F}$ is a family of graphs containing some disconnected graph H, then Planar-$\mathcal{F}$-Deletion has not FII (in general).

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- Then, if we take $1 \leqslant n < m$,

$$\begin{aligned} \pi(G_n \oplus H_{n-1}) - \pi(G_m \oplus H_{n-1}) &= (n-1) - (n-1) = 0, \\ \pi(G_n \oplus H_m) - \pi(G_m \oplus H_m) &= n - m < 0. \end{aligned}$$

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- Thus, $G_n, G_m \notin$ same equiv. class of $\sim_{\Pi,1}$ whenever $1 \leqslant n < m$.

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Lemma (The parameter does not increase)

$\forall$ fixed t, $\exists$ finite set $\mathcal{R}_t$ of t-boundaried graphs s.t. for each t-boundaried graph $G \in \mathcal{G}_t$ $\exists$ $G' \in \mathcal{R}_t$ s.t. $G \equiv_{\Pi,t} G'$ and $\Delta_{\Pi,t}(G, G') \geqslant 0$.

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Lemma (Replacing protrusions of constant size)

For $t \in \mathbb{N}$, suppose that the set $\mathcal{R}_t$ of representatives of $\equiv_{\Pi,t}$ is given. If W is a t-protrusion of size at most a fixed constant c, then one can decide in constant time which $G' \in \mathcal{R}_t$ satisfies $G' \equiv_{\Pi,t} G[W]$.

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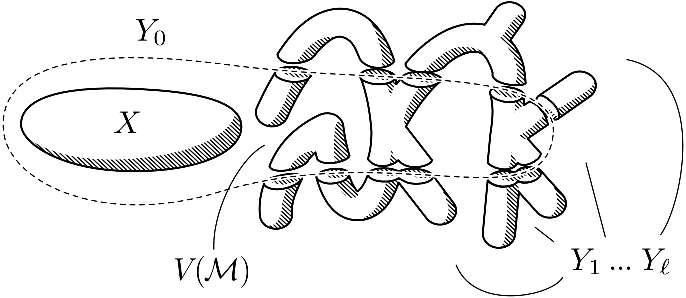
It runs in polynomial time... given the sets of representatives!

Protrusion decompositions (in case someone forgot!)

An (α, t)-protrusion decomposition of a graph G is a partition $\mathcal{P} = Y_0 \uplus Y_1 \uplus \cdots \uplus Y_\ell$ of $V(G)$ such that:

- for every $1 \leqslant i \leqslant \ell$, $N(Y_i) \subseteq Y_0$;
- for every $1 \leqslant i \leqslant \ell$, $Y_i \cup N_{Y_0}(Y_i)$ is a t-protrusion of G;
- $\max\{\ell, |Y_0|\} \leqslant \alpha$.



(Figure by Felix Reidl)

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   1. $|Y_0| = O(k)$;
   2. $|Y_i| \leqslant \rho'_{\Pi}(2t + \omega_{\mathcal{H}})$ for $1 \leqslant i \leqslant \ell$; and
   3. $\ell = O(k)$.

Next subsection is...

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- Constructing the kernels? Finding the sets of representatives!!

Limits of our approach and further research

- For which notions of sparseness (beyond H-topological-minor-free graphs) can we use our technique to obtain polynomial kernels?

  1. A class $\mathcal{G}$ of graphs locally excludes a minor if $\forall r \in \mathbb{N}$, $\exists H_r$ s.t. the r-neighborhood of a vertex of any graph of $\mathcal{G}$ excludes H_r as a minor.

     (includes H-minor-free but incomparable with H-topological-minor-free)

     Except for a very restricted case, our technique fails.

  2. Graphs of bounded expansion (contains H-topological-minor-free)?

     Obtaining a kernel for Treewidth-t Vertex Deletion on graphs of bounded expansion is as hard as on general graphs.

     Best known kernel: $k^{O(t)}$. [Fomin, Lokshtanov, Misra, Saurabh '12]

- Constructing the kernels? Finding the sets of representatives!!
- Explicit constants? Lower bounds on their size?

Gràcies!