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Efficient Method for Magnitude Comparison in RNS Based on Two Pairs of Conjugate Moduli

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Outline



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1. Motivation
2. Class of Moduli Sets
3. Method for comparing magnitude in RNS
4. Typical application: RNS motion estimator
5. Conclusions and future work



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- Carry free arithmetic
 - Residue Number Systems (RNS) allows to parallelize +, -, *
- No general efficient method for comparison in RNS
 - to convert from residues to positional code: CRT requires modulo M operations and MRC is a sequential method!
 - [Miller 86, Dimauro 93, Wang 99]: computationally demanding, not suitable for hardware implementation
- To propose a method for comparing RNS numbers considering a representative class of moduli sets.



- New class of multi-moduli sets that rely on pairs of conjugate moduli:
$$S = \{m_1, m_1^*, \dots, m_k, m_k^*\} = \{2^{n_1}-1, 2^{n_1}+1, \dots, 2^{n_k}-1, 2^{n_k}+1\}$$
 - Only Mersenne rings and Fermat rings
- S is not a set of pairwise relatively prime, however modified CRT [Wang98] allows to obtain an integer from residues
- This class of multi-moduli leads to two-level residue number systems
- Important sub-class S' : two pairs of balanced conjugate moduli sets

$$S' = \{m_1, m_1^*, m_2, m_2^*\} = \{2^n-1, 2^n+1, 2^{n+1}-1, 2^{n+1}+1\}$$



- For each of the two level we can use the modified CRT to compute RNS-to-binary

$$X = X_1 + m_1 \left\langle \left(\frac{m_1}{d} \right)^{-1} \frac{X_2 - X_1}{d} \right\rangle_{\frac{m_2}{d}}$$

– $d = \text{GCD}(m_1, m_2) \Rightarrow$ pairwise relatively prime $\text{CRTIII} \equiv \text{MRC}$

$$d = 1 \rightarrow X_1 = x_1^* + (2^n + 1) \left\langle 2^{n-1} (x_1 - x_1^*) \right\rangle_{2^{n-1}}$$

$$d = 3 \rightarrow X = X_2 + \left\langle \left(\frac{2^{2(n+1)} - 1}{3} \right)^{-1} \frac{X_1 - X_2}{3} \right\rangle_{\frac{2^{2n-1}}{3}}$$

- Very important for us is that the range is odd

$$M = \frac{(2^{2n} - 1)(2^{2n+2} - 1)}{3}$$



- Unsigned integer numbers (A, B) can be compared by subtraction:

$$C = \begin{cases} A - B & \text{for } A \geq B \\ M - A - B & \text{for } A < B \end{cases}$$

- Based on the well known mathematical axiom:
 - *the subtraction of two numbers with the same parity leads to an even number and the subtraction of two numbers with different parities leads to an odd number*
- and taking advantage that **M is odd** we can answer the question :
 - is $A \geq B$ or not?

- PREPOSITIONS

- $A \geq B$ iff:
 - A and B have the same parity and C is an even number
 - A and B have different parities but C is an odd number.
- $A < B$ iff:
 - A and B have the same parity and C is an odd number
 - A and B have different parities but C is an even number.

So we have to compute the parity of A, B and C!



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- Problem: how to directly compute the parity of a RNS number?
 - without computing the number back to a traditional weighted system!
- The parity of an integer X in the range [0, M-1] represented on the $\{2^n-1, 2^n+1, 2^{n+1}-1, 2^{n+1}+1\}$ moduli set can be computed by:
$$\langle X \rangle_2 = \left\langle \left\langle X_2 \right\rangle_2 \oplus \left\langle \left\langle X_1 - X_2 \right\rangle_{2^{2n-1}} \right\rangle_2 \right\rangle_2$$
- by converting X1 and X2 in the 1st-level we also just need shift and one's complement addition

$$X_1 = x_1^* + (2^n + 1) \times \left\langle 2^{n-1} (x_1 - x_1^*) \right\rangle_{2^n-1}$$

$$X_2 = x_2^* + (2^{n+1} + 1) \times \left\langle (2^n (x_2 - x_2^*)) \right\rangle_{2^{n+1}-1}$$



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Algorithm 1 Comparison of the numbers A, B represented in RNS $(a_1, a_1^*, a_2, a_2^*, b_1, b_1^*, b_2, b_2^*)$.

- 1: $c_1 = \langle a_1 - b_1 \rangle_{2^n-1}; c_1^* = \langle a_1^* - b_1^* \rangle_{2^n+1};$
 $c_2 = \langle a_2 - b_2 \rangle_{2^{n+1}-1}; c_2^* = \langle a_2^* - b_2^* \rangle_{2^{n+1}+1};$
- 2: $(A_1, A_2) = \text{1st-level-converter}(a_1, a_1^*, a_2, a_2^*); \quad \{(15) \text{ and } (16)\}$
 $(B_1, B_2) = \text{1st-level-converter}(b_1, b_1^*, b_2, b_2^*); \quad \{(15) \text{ and } (16)\}$
 $(C_1, C_2) = \text{1st-level-converter}(c_1, c_1^*, c_2, c_2^*); \quad \{(15) \text{ and } (16)\}$
- 3: $\overline{P_A} = LSB(\langle A_1 - A_2 \rangle_{2^{2n}-1}) \oplus LSB(A2); \quad \{'1' \text{ if } X \text{ even}\}$
 $\overline{P_B} = LSB(\langle B_1 - B_2 \rangle_{2^{2n}-1}) \oplus LSB(B2);$
 $\overline{P_C} = LSB(\langle C_1 - C_2 \rangle_{2^{2n}-1}) \oplus LSB(C2);$
- 4: if $P_A \oplus P_B \oplus P_C = '1'$ then
- 5: $A \geq B$ is TRUE;
- 6: else
- 7: $A < B$ is TRUE;
- 8: end if

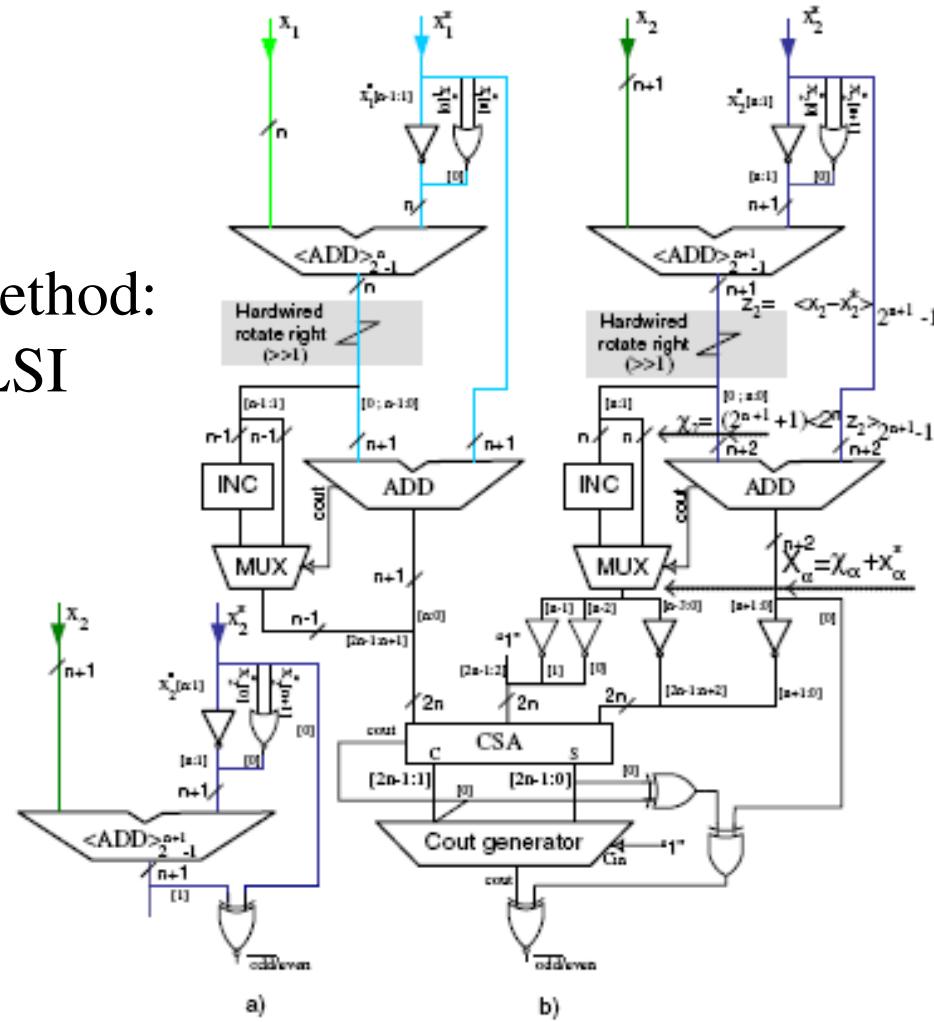


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Parity detection method:
suitable for VLSI



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Maximum Operation Size (MOS)		
Algorithm	MOS (n,M)	MOS (n=4)
Miller	$\approx 4M$	≈ 86955
Dimauro	modulo($\approx 2^{3n}+2$)	modulo(≈ 4098)
Wang	modulo($2^{2n}-1$)	modulo(255)
Proposed	modulo($2^{n+1}+1$)	modulo(33)



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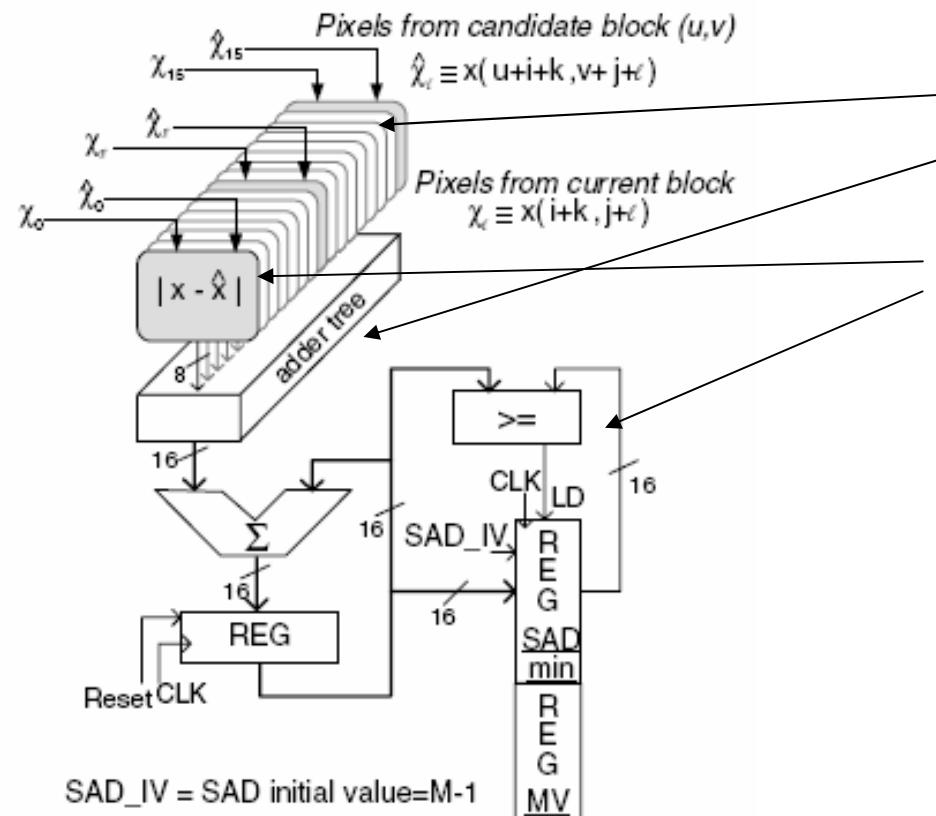
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Typical application: RNS motion estimator

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- Tradicional architecture



Subtractors and
Adders



Comparators:
proposed hardware

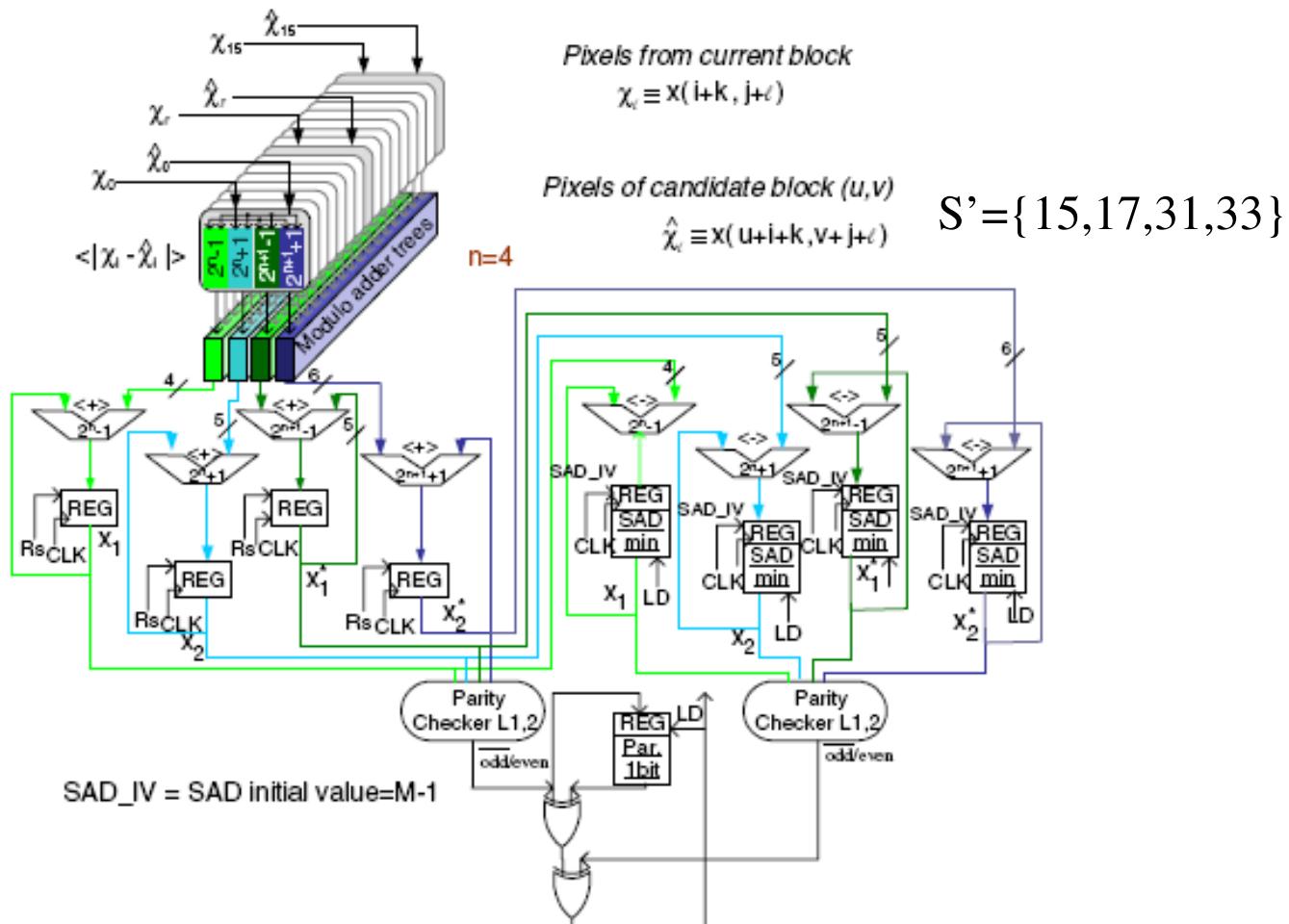


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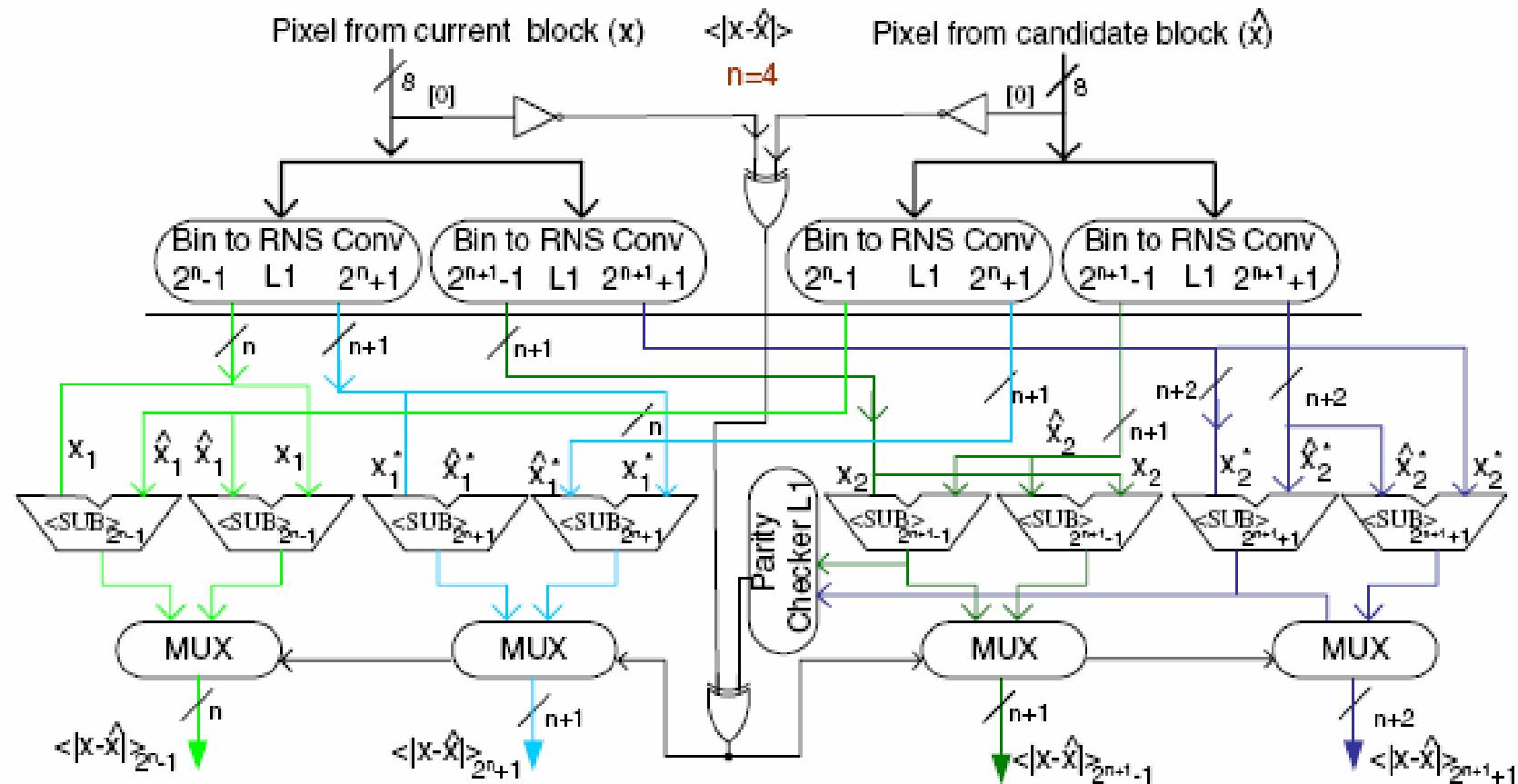
RNS motion estimator

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Experimental Results: RNS motion estimator



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- SAD unit implemented in a FPGA with arithmetic units directly mapped on Look-Up-Tables (LUT)
- FPGA Xilinx VirtexII Pro (xc2vp50-7)
- Synthesis with ISE (8.2) tools

Slices (% total)	BRAMs (% total)	Freq. MHz	Latency Cycles	Throughput Blocks/s
246 (1%)	211 (90%)	254	12	1.5×10^7



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Conclusions and future work



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- New efficient method is proposed for magnitude comparison in RNS based on two pairs of conjugate moduli
- This is the first method leading to VLSI architectures with practical interest for comparing the magnitude of numbers in RNS
 - Efficient RNS minimum SAD unit was already implemented in FPGA
 - We are implementing a SAD unit on an ASIC (0.18µm CMOS)
- We are now extending the idea to other moduli sets, all with a common characteristic: **M odd**



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