

ORCCAD, a framework for safe robot control design and implementation

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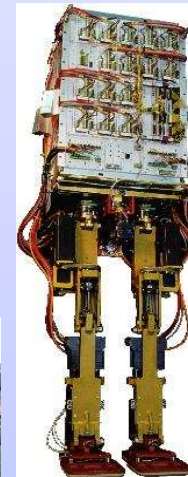
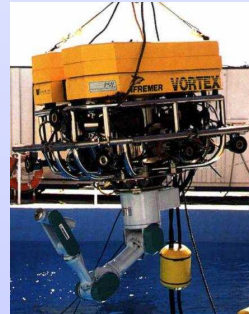
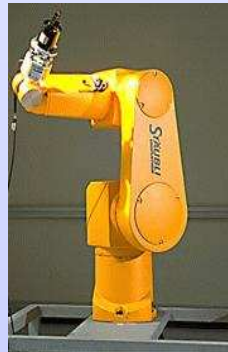
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<http://sed.inrialpes.fr/Orccad/>

Open Robot Controller Computer Aided Design

- ORCCAD : <http://sed.inrialpes.fr/Orccad/>
- A set of concepts, methods and programming tools for robotics
 - Handling automatic control and reactivity
 - Tools and methods used from design to run-time



General framework

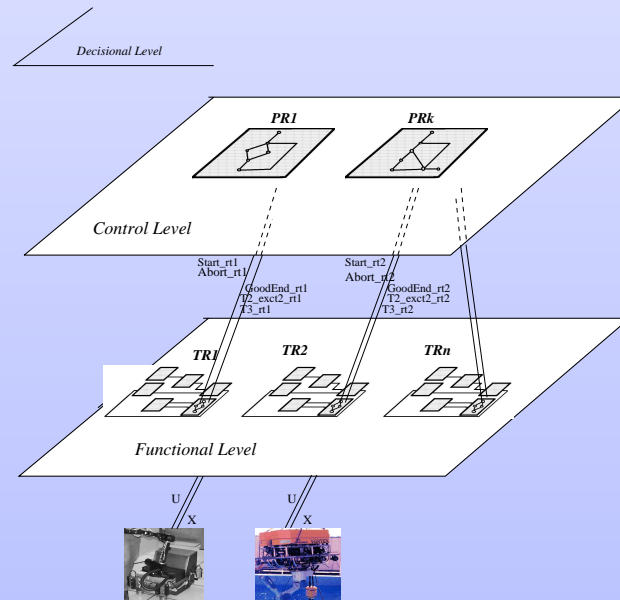
- Complex mechanical systems: many degrees of freedom, non-linearities, coupling
$$\Gamma = M(q)\ddot{q} + N(q, \dot{q})$$
- Possible switches in the model (e.g. walking machines, dynamic vision)
- Interaction with the environment, the environment is dynamic and uncertain
- Needs for reliability (e.g. hostile environment, cooperation with humans)
- Hybrid aspects
 - Sampling for sensing and periodic servo-loops
 - Discrete Events for decisional processes

Guidelines

- Many robotic actions can be stated efficiently in the framework of Automatic Control theory
Task-function approach, Sensor-based control...
(*Bottom-up approach*)
- A robotic system has several kinds of users, each of them must be provided with dedicated CAD tools
(*Different use-cases...*)
- Software for robotics must be reliable and reusable
(*Validation, formal verification, automatic code generation...*)
- The overall system performance strongly relies on implementation of real-time programs
(*Various constraints and co-design*)

Main structures

- Robot-Tasks (RT) model basic robot's actions
Control law + Logical behaviour
Control aspects are predominant
- Robot-Procedures (RP): logical and temporal composition of RTs and RPs for incremental design of actions with variable complexity
Logical aspects are predominant



the Robot-Task: an hybrid object

- Algorithmic aspects:
Parameterized control law the structure of which is invariant along the RT duration
 - ⇒ Trees of object classes
 - ⇒ the leaves are real-time tasks
 - ⇒ Instantiated via a G.U.I.
- Reactive aspects:
Specification of the logical behaviour of the RT
 - ⇒ Encoded with Esterel
 - ⇒ Coherency with higher levels (RPs)
 - ⇒ Formal verification

The reactive behaviour encapsulates numerical computations

The RT external view is a Discrete Event System.

Control law design

Task-function : control goal in the task space

Control law structure:

$$\Gamma = -k\hat{M} \left(\frac{\widehat{\partial e}}{\partial q} \right)^{-1} G \left(\mu D e + \frac{\widehat{\partial e}}{\partial q} \dot{q} + \frac{\widehat{\partial e}}{\partial t} \right) + \hat{N} - \hat{M} \left(\frac{\widehat{\partial e}}{\partial q} \right)^{-1} \hat{f}$$

Several control laws can be given according to the choice of models, e.g.

- \hat{M} diagonal and $\hat{N} = 0$: fixed gain decentralized PD
- explicit computation of \hat{M} and \hat{N} : computed torque control...

Modular Specification of a Robot-Task

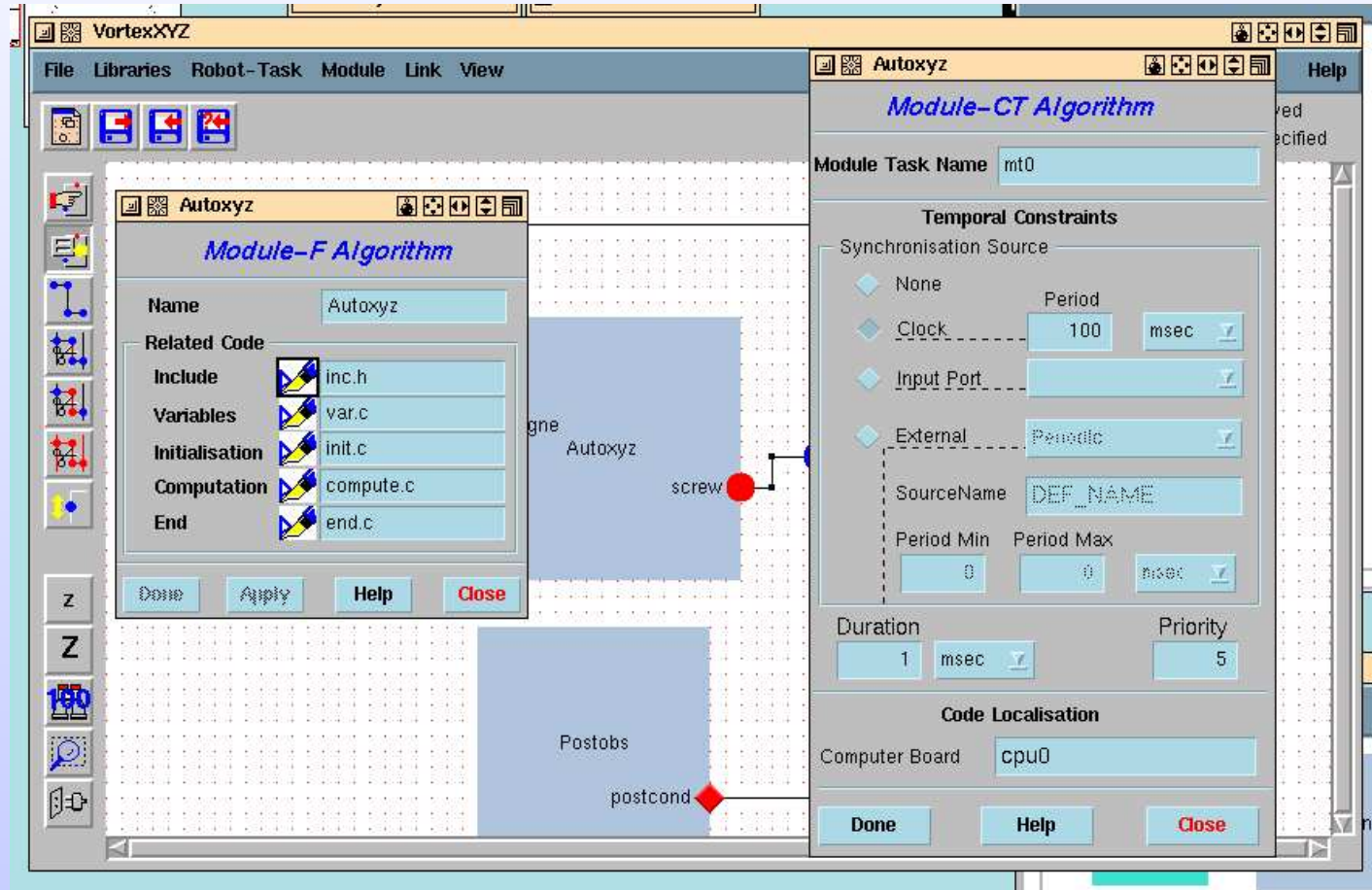
The screenshot displays the Orccad software interface for modular specification. The main workspace shows a block diagram with modules: Excavator (orange), Kin (blue), Control (grey), and Exc_Atr (teal). Connections are shown between these modules, including a 'Gamma' port and a 'q' port. A 'Std Module' palette on the right lists available modules: Excavator, Kin, Exc_Atr, Gene_Traj_SE3, Man_Atr, Man_Cont, Man_Traj, Man_PhR, and Simple_Cont. Two dialog boxes are open: 'Kin' (Module-F Algorithm) and 'Event Port' (End_traj). The 'Kin' dialog shows related code files: inc.h, var.c, init.c, compute.c, and end.c. The 'Event Port' dialog shows configuration for an 'End_traj' event, including Name, Type (Post-Condition), Timeout (1.00 sec), and Predicate (End_traj).

Modules

Several kinds of *Modules*

- class *algorithms*:
 - numerical computation and/or observers
 - temporal attributes : period, duration, synchronizations...
- class *physical resource*: interface with the process (drivers)
- class *Robot-task Automaton*: reactive behaviour (e.g. encoded with ESTEREL)

Algorithmic module



- Functional attributes: name and location of C files
- Temporal attributes : Synchro., Priority, CPU ID, WCET

Control and temporal constraints

Goal : achieve control performance w.r.t. available computing power

Theoretic background : few quantitative results, esp. for non-linear systems

Current practice :

- sampling rate : $0.15 < \omega h < 0.5$ (rule of thumb)
- latency \ll period, or change the control algorithm
- jitter ???, very sensitive on identification

Specify time dependencies through synchronisation

Assign priorities according to the relative urgency of tasks

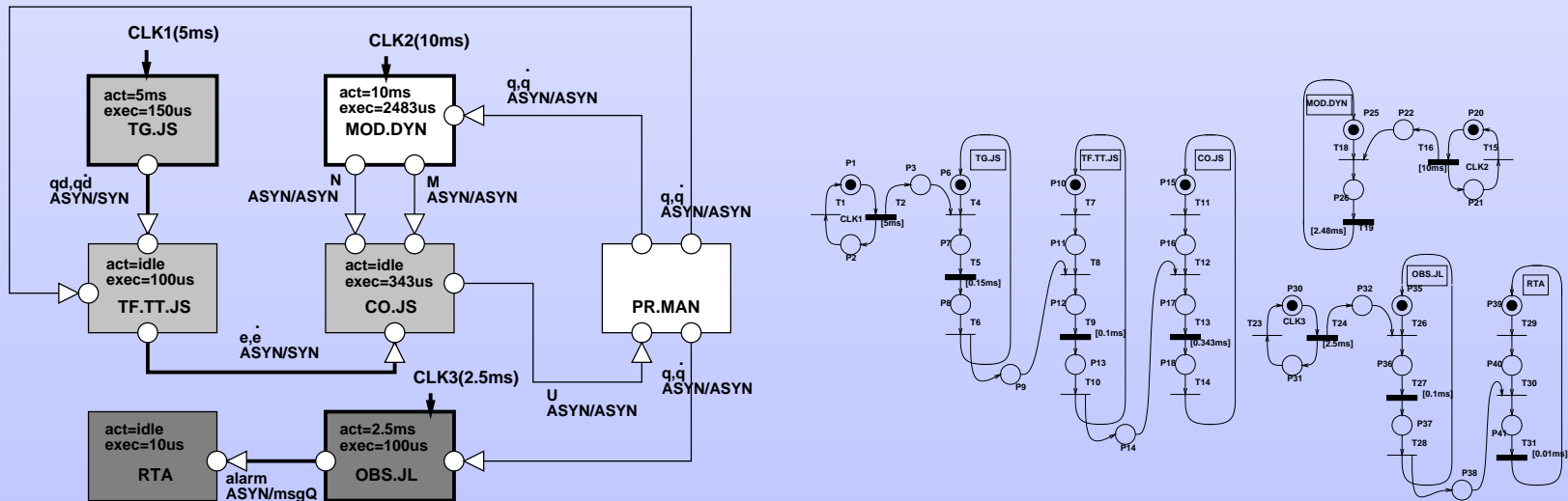
Synchronised tasks must have the same priority

The best data is the freshest one \Rightarrow single slot links

Take care with shared data and synchronization side-effects

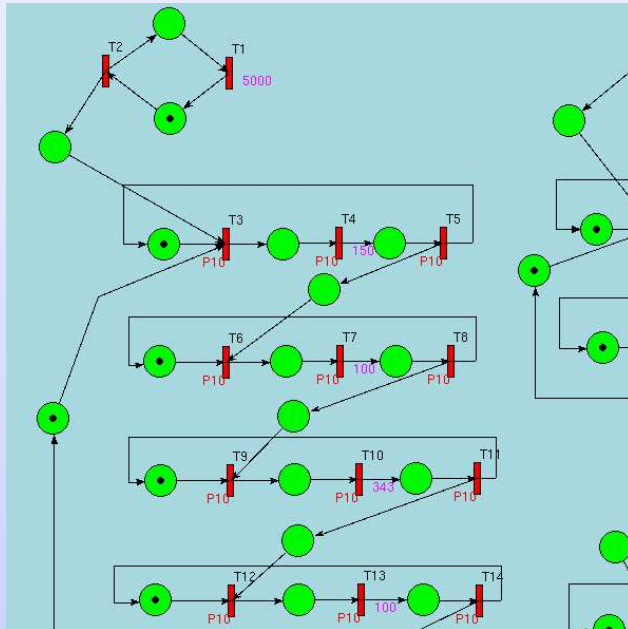
Verification of temporal constraints

- Partial synchronization to specify precedence constraints
- Timed Petri nets (marked graphs)
- Deadlock analysis by checking the liveness of p-invariants
- Analysis of temporal properties using linear models in the $(\max, +)$ algebra (periodicity, stability, respect of deadlines...)



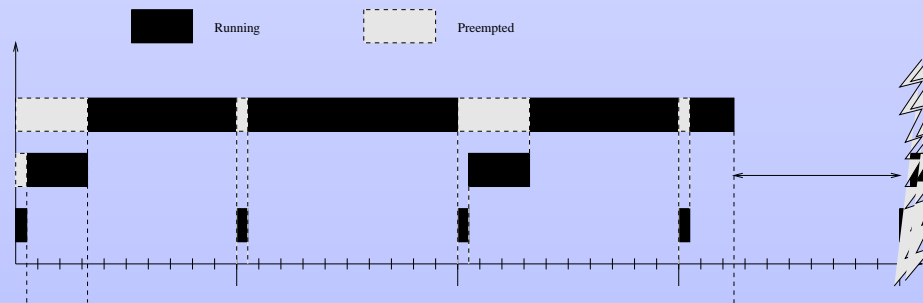
Schedulability analysis

complex synchronisation and scheduling schemes :
 ORCCAD GUI => ERS solver in (max,plus) algebra

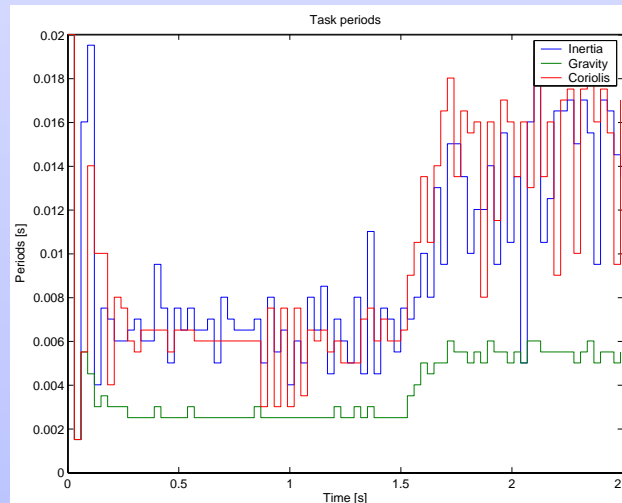
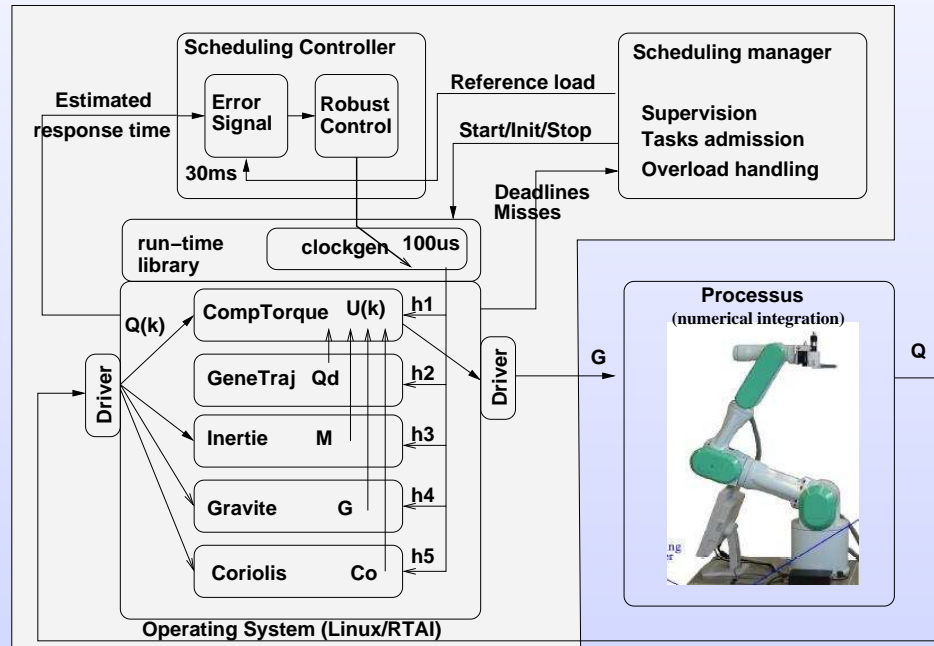


```

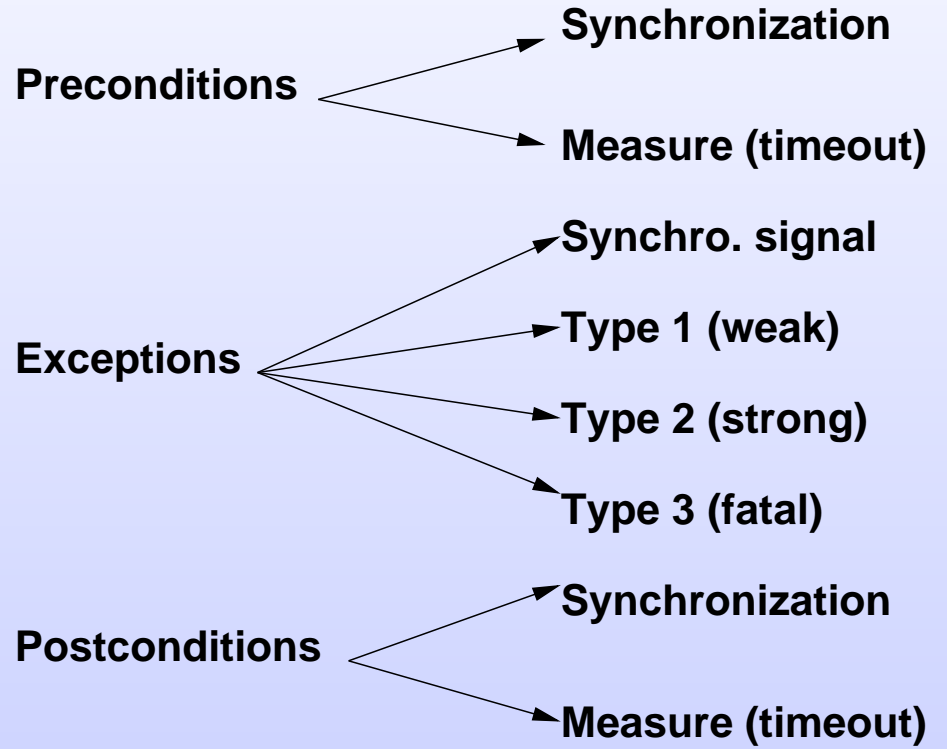
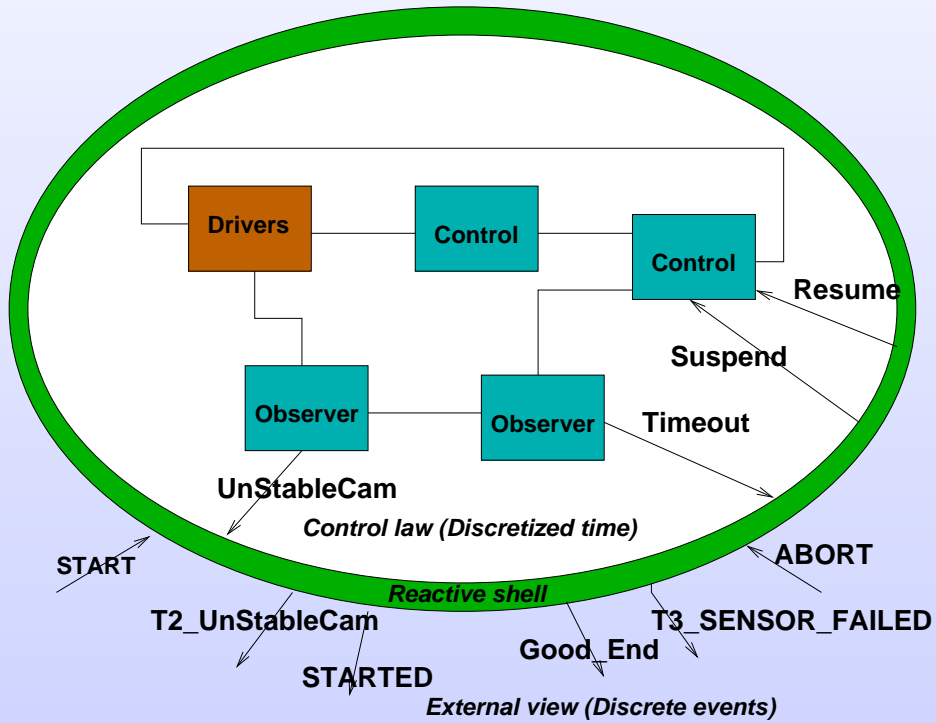
Période du système (temps contracté): 8174/1
Période du système (temps réel): 10000/1
DBG: activité autonome (temps contracté):
[[[0,6280]]
DBG: activité préemptée:
1(8106)[0(1894)1(8106)]
Taux d'inactivité: 1894/10000
DBG: activité préemptée (normalisée):
1(8106)[0(1894)1(8106)]
    
```



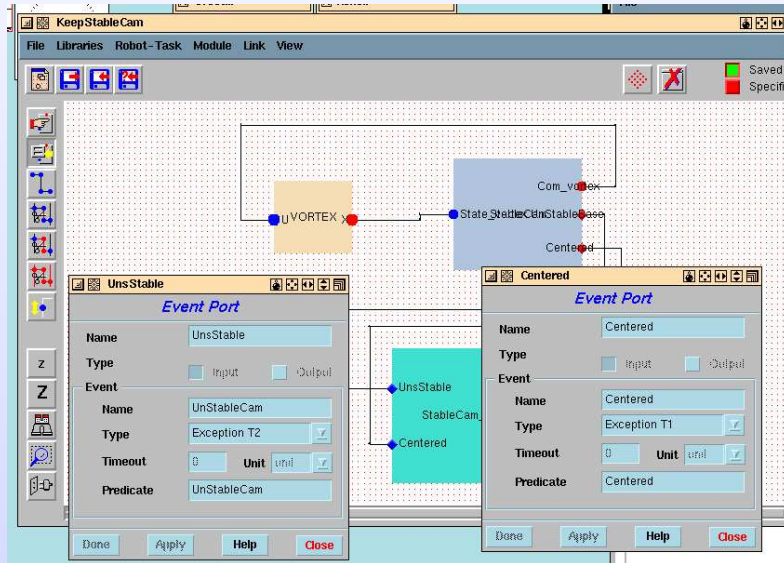
Feedback scheduling



Events



The RTA module



```

trap T2 in
trap BF in
[
  await Abort_Local_KeepStableCam;
  emit StartTransite_VORTEX;
  emit KeepStableCamTransite(?KeepStableCam_Start);
  exit Abort;
]
||
loop
  await Centered; emit TreatCentered(?KeepStableCam_Start)
end loop
||
await UnStableCam; emit T2_UnStableCam; exit T2
||
[[[
halt;
]]]; exit BF; ]
    
```


Esterel

- Synchronous language with imperative style
- Modularity using parallelism at specification time
- Powerfull escape mechanisms
- Multiform time
- Deterministic behaviour \Rightarrow formal verification
- Inputs and outputs are signals (pure or valued)
- Communication based on synchronous diffusion
- Compiles into efficient deterministic automata

Main statements

emit S

await S

p ; q

p || q

call P (X1,X2)(e1,e2)

exec T ()() return R

loop p end

loop p each S end

present S then p else q end

abort p when S

weak abort p when S

diffusion of a signal

await the next occurrence

sequence

parallel

extern procedure call

extern task launch

infinite loop

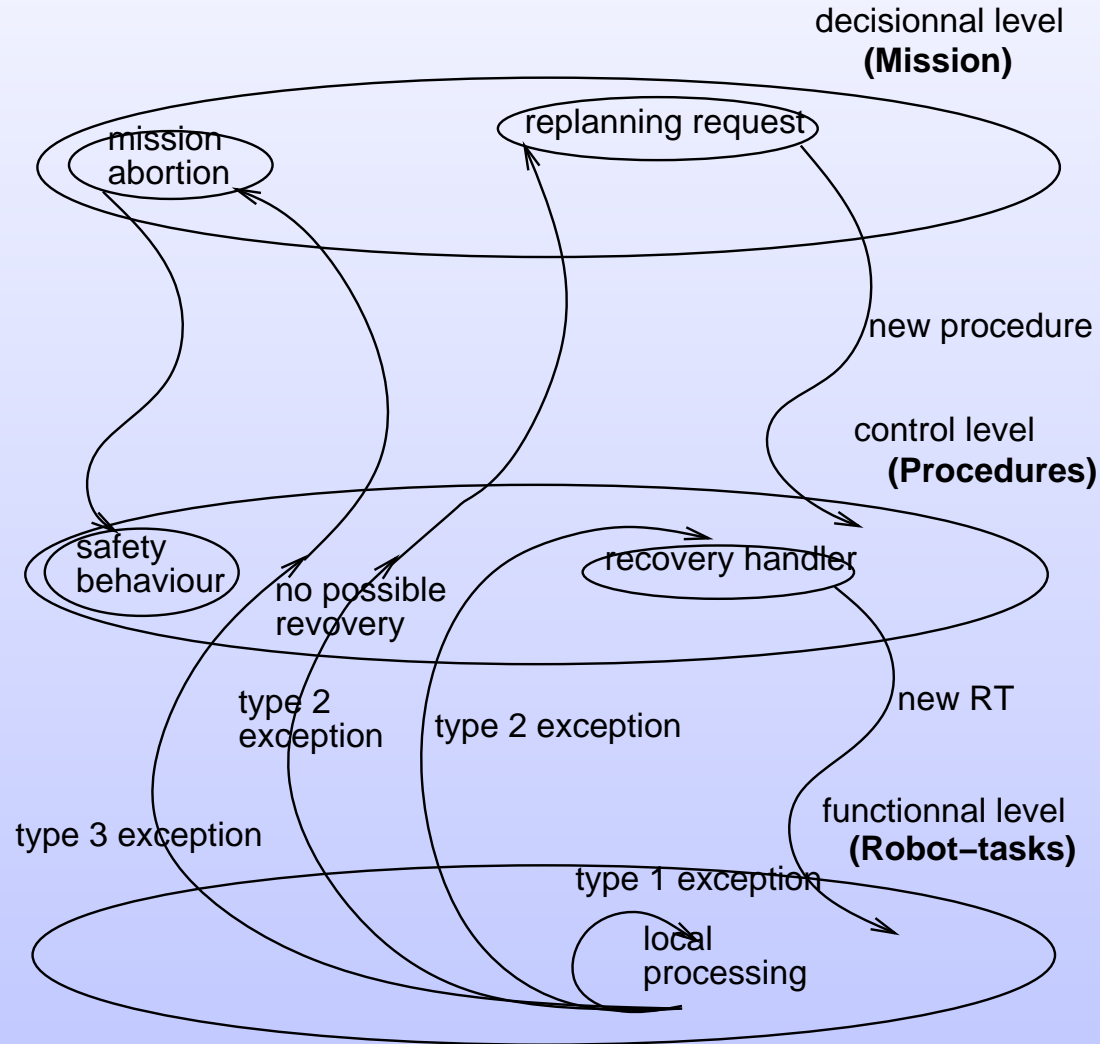
temporal loop

branching

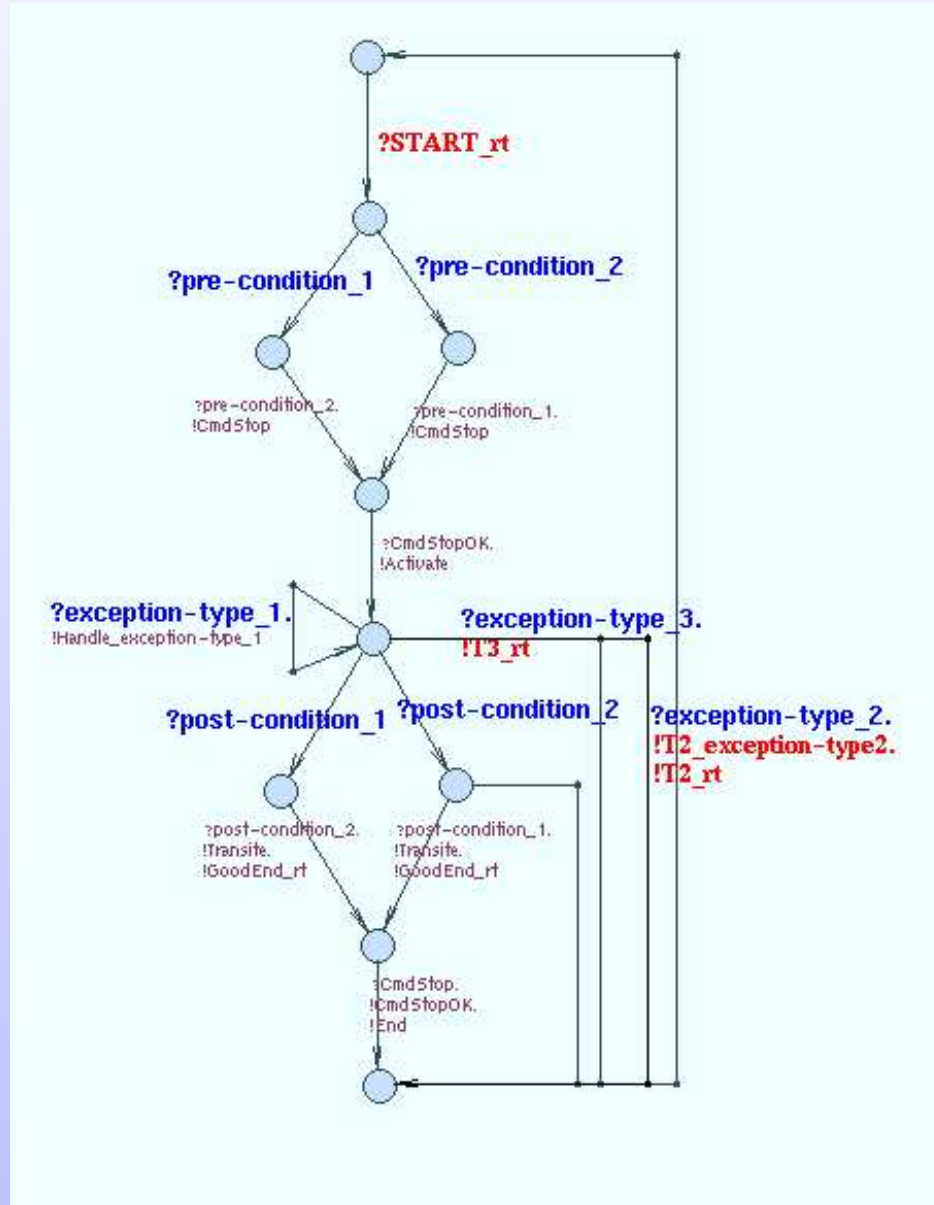
strong preemption

weak preemption

Exception processing



A Robot-Task Controller



Robot-Procedure specification

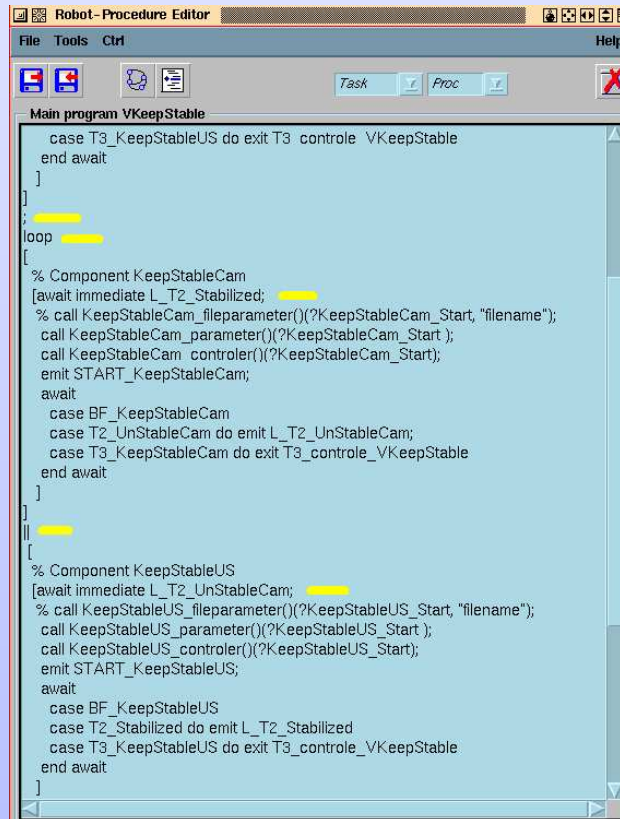
Incremental design of complex actions (nominal and degraded modes)

Exception handling

Semi-automatic code generation

ESTEREL : sequence, parallel, loops, watchdogs, traps...

We don't have a FSM in mind at design time



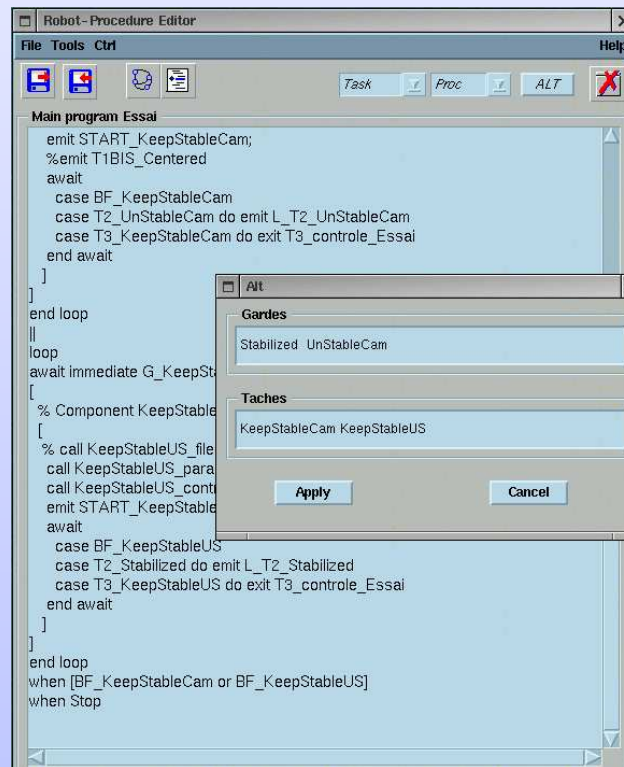
```

Robot-Procedure Editor
File Tools Ctrl Help
Task Proc
Main program VKeepStable
case T3_KeepStableUS do exit T3_controle_VKeepStable
end await
]
]
]
loop
[
% Component KeepStableCam
[await immediate L_T2_Stabilized;
% call KeepStableCam_fileparameter()(?KeepStableCam_Start, "filename");
call KeepStableCam_parameter()(?KeepStableCam_Start);
call KeepStableCam_controler()(?KeepStableCam_Start);
emit START_KeepStableCam;
await
case BF_KeepStableCam
case T2_UnStableCam do emit L_T2_UnStableCam;
case T3_KeepStableCam do exit T3_controle_VKeepStable
end await
]
]
||
[
% Component KeepStableUS
[await immediate L_T2_UnStableCam;
% call KeepStableUS_fileparameter()(?KeepStableUS_Start, "filename");
call KeepStableUS_parameter()(?KeepStableUS_Start);
call KeepStableUS_controler()(?KeepStableUS_Start);
emit START_KeepStableUS;
await
case BF_KeepStableUS
case T2_Stabilized do emit L_T2_Stabilized
case T3_KeepStableUS do exit T3_controle_VKeepStable
end await
]
]
]

```

Specification of a guarded alternance

- actions are guarded by T2 emitted by other actions
- actions are mutually exclusive
- no dead-lock nor endless loop



The screenshot shows the 'Robot-Procedure Editor' window. The main text area contains the following code:

```

Main program Essai
emit START_KeepStableCam;
%emit T1BIS_Centered
await
case BF_KeepStableCam
case T2_UnStableCam do emit L_T2_UnStableCam
case T3_KeepStableCam do exit T3_controle_Essai
end await
]
end loop
||
loop
await immediate G_KeepSt
[
% Component KeepStable
[
% call KeepStableUS_file
call KeepStableUS_para
call KeepStableUS_cont
emit START_KeepStable
await
case BF_KeepStableUS
case T2_Stabilized do emit L_T2_Stabilized
case T3_KeepStableUS do exit T3_controle_Essai
end await
]
]
end loop
when [BF_KeepStableCam or BF_KeepStableUS]
when Stop
  
```

An 'Alt' dialog box is open over the code, with the following fields:

- Guardes:** Stabilized UnStableCam
- Taches:** KeepStableCam KeepStableUS

The dialog box has 'Apply' and 'Cancel' buttons at the bottom.

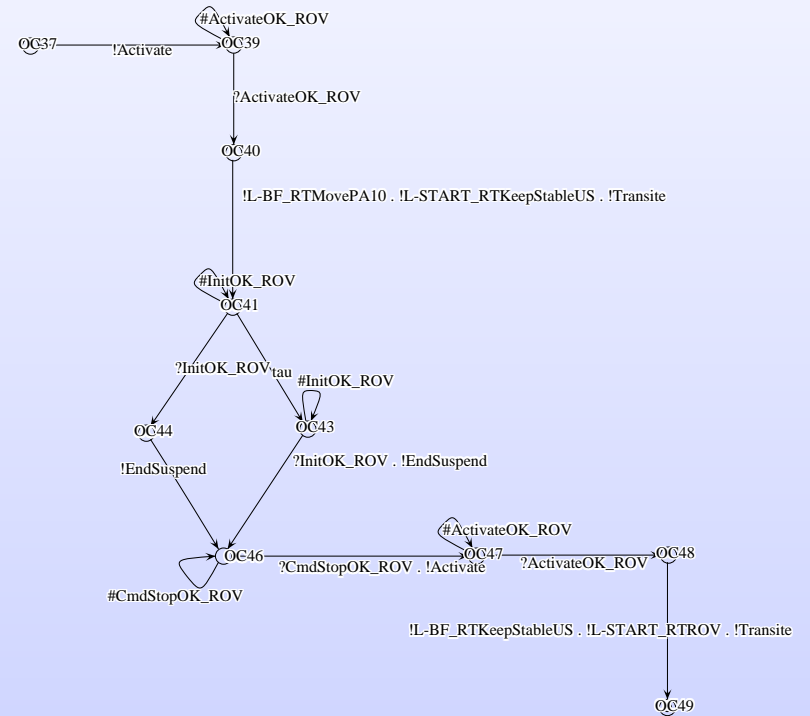
Maestro : a domain specific language

```
do
  SEQ(do
    KeepStableUS
    until Stabilized
  ;
  loop
    PAR(when T2_Stabilized when T2_UnstableCam
      do
        KeepStableCam
        until UnStableCam
      do
        KeepStableUS
        until Stabilized)
    end loop
  )
until Stop
```

a graphical version would be better???

Switching RTs

Robot Task 1	Robot Task 2	Switching signals from/to automaton	Comments
•		→ post-condition or type 2 exception	Decision for switching to RT2
←		• Transite	
	←	• Resume	Resume real-time tasks (MTs)
	•	→ Resume_ok	
	←	• Init	Task initialization Checking for pre-conditions
	•	→ Init_ok	
←		• Control_stop	Stop of control application
•		→ Control_stop_ok	
	←	• Control_start	Start of control application
←		• Suspend	
•		→ Suspend_ok	



Complex logic + real-time scheduling (overlap during transite)

Formal verification : requirements

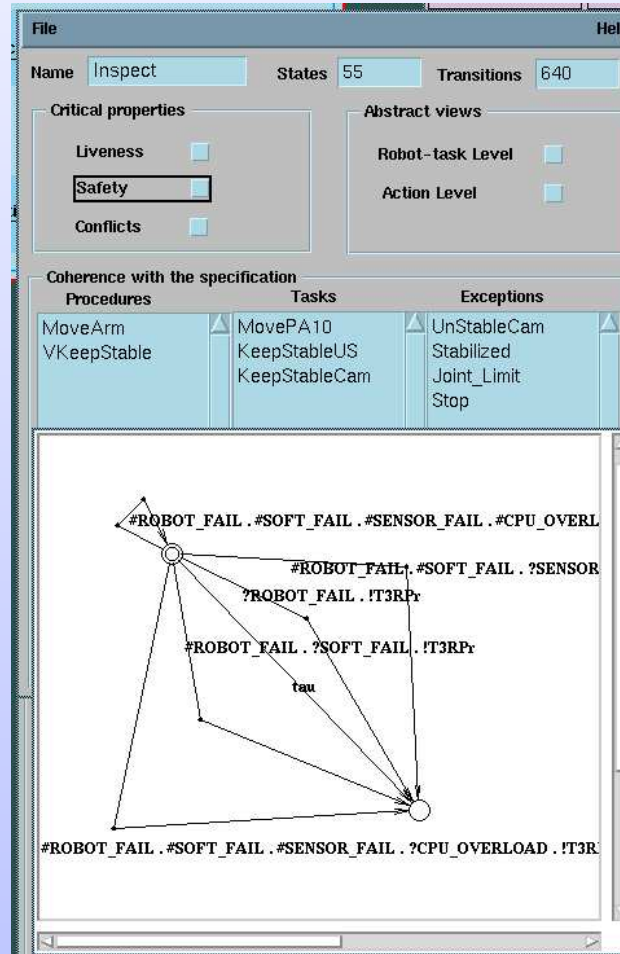
Once a mission has been defined, check that:

- its specification is correct (corresponds to the desired goal)
 - its programming conforms to specification
 - real-time implementation does not disturb its behaviour
-
- Safety properties (fatal exceptions are always correctly handled)
 - Liveliness properties (the goal is always reached in nominal executions)
 - Conflicts detection (mutual exclusion)
 - Conformity with the requirements
 - Help to specification (abstract views)

Safety properties

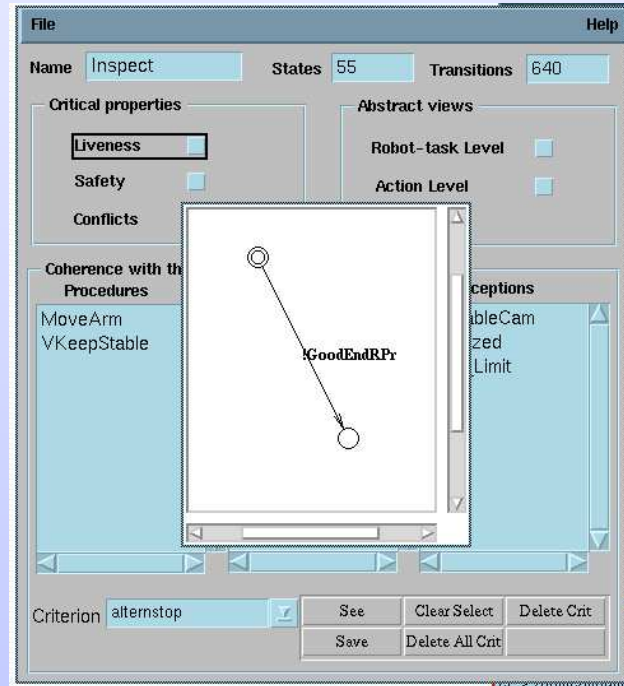
fatal exceptions are always correctly handled

abstract action: “*Error = /Water_Leak? and not /Accent!*”



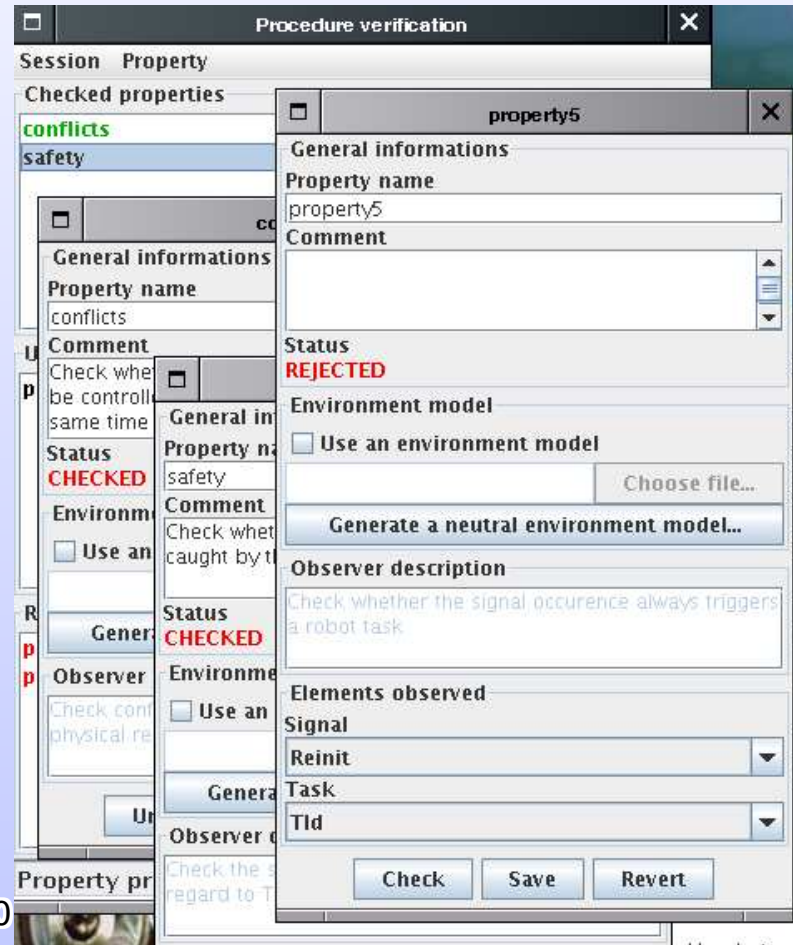
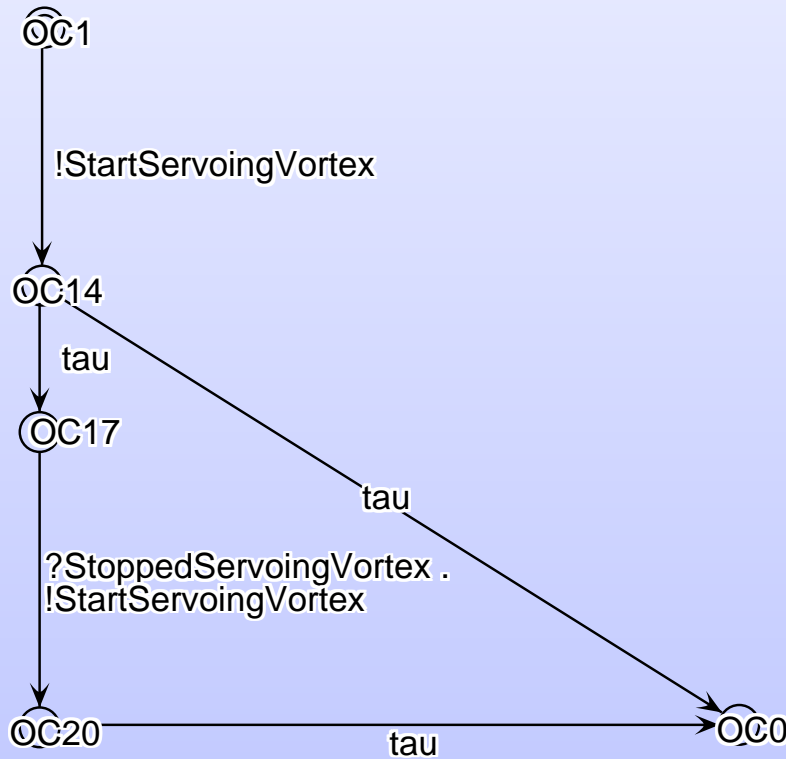
Liveliness properties

the goal is always reached in nominal executions
e.g. there are no endless loops...



Conflicts detection

two control laws must not compete to control the same robot



Conformity with the requirements

Start KeepStable;

weak abort

Start KeepStableUS;

[await T2_Stabilized

Start KeepStableCam

await T2_Unstable

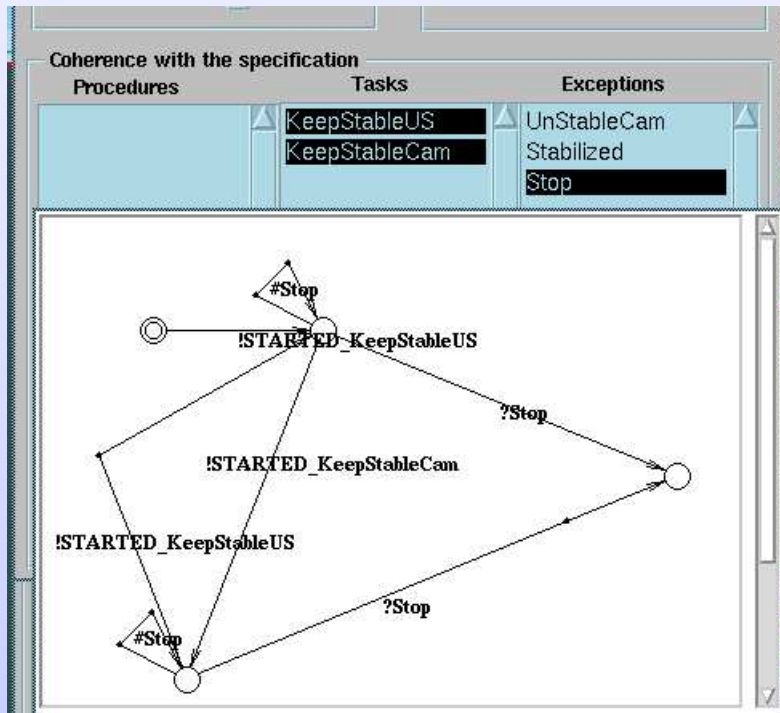
||

await T2_Unstable;

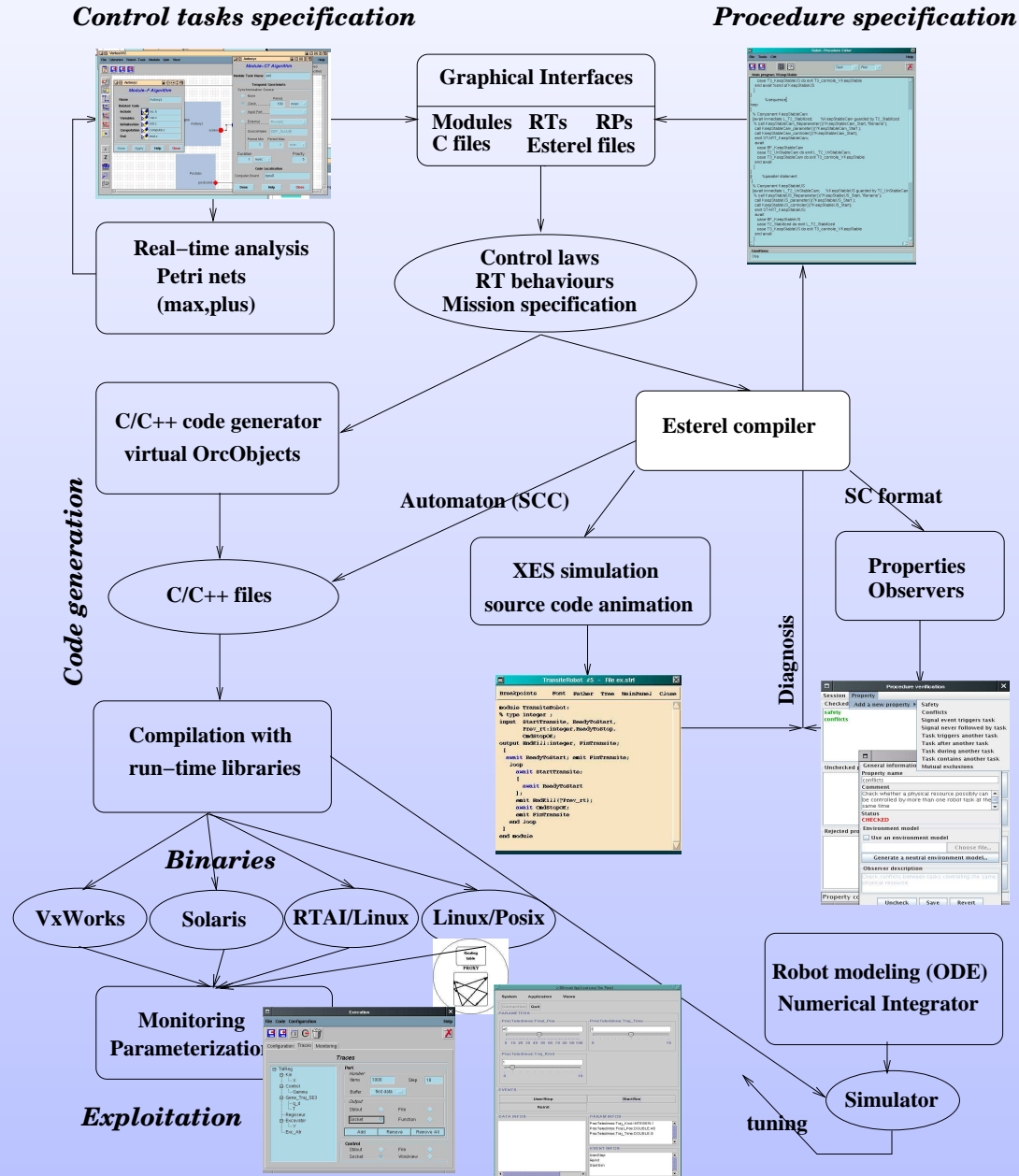
Start KeepStableUS

await T2_Stabilized

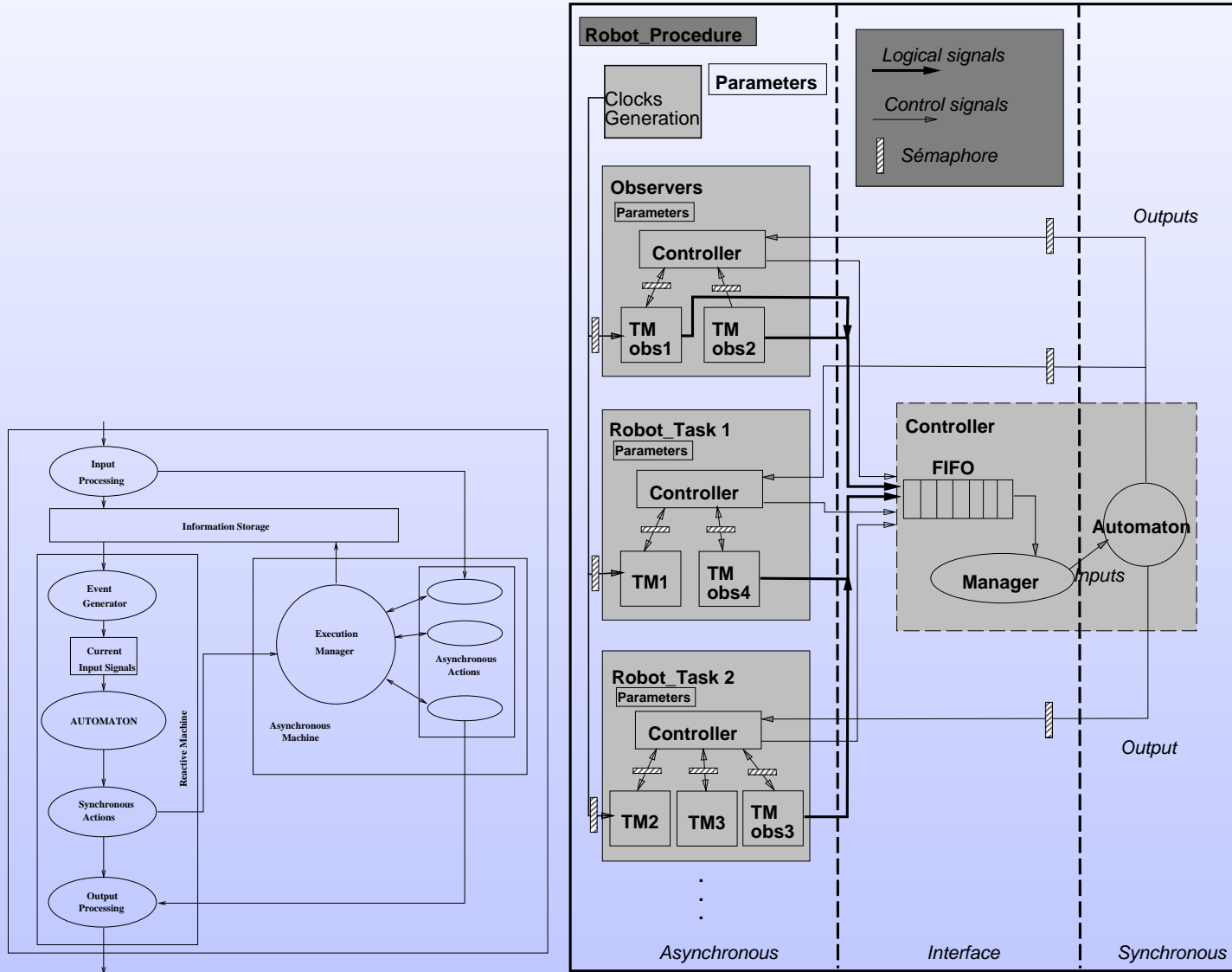
when Stop



Summary of Orccad tools



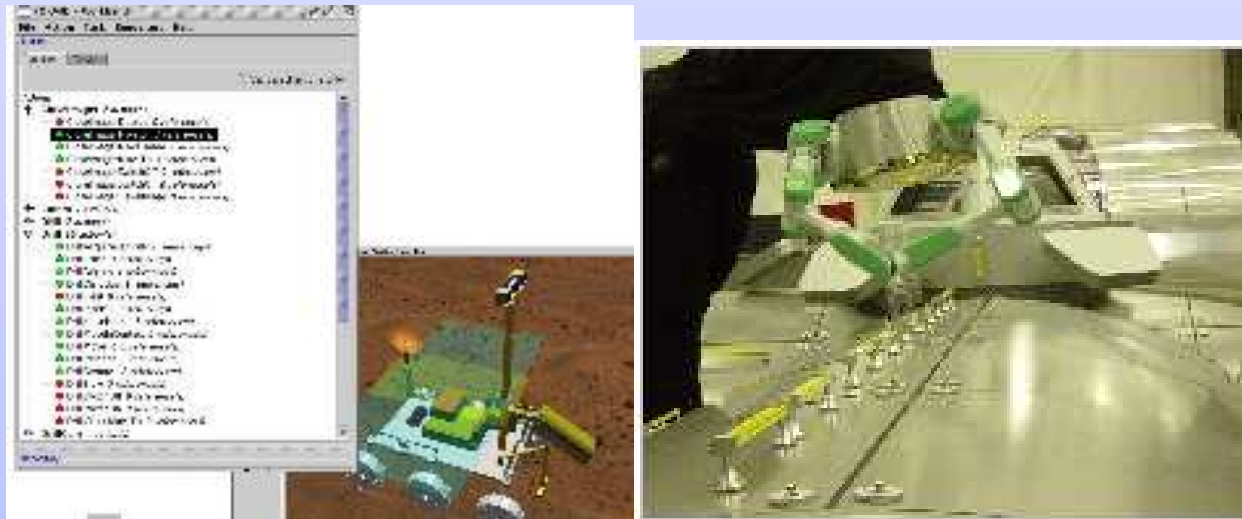
Real-time implementation



only uses the basic features of the RTOS

Using Orccad tools

- BIP2000 designers : control and/or mechanical engineering
- flexible and feedback scheduling
- underwater arm/vehicle high-level coordination
- specification/simulation of missions for a Mars rover
- control of vision-based tasks (Vimanco)
- teleoperation interface with force-feedback (Teleman)



Current release : Orccad V3.1

- Edition of modules : algorithmic, automaton and PhR
- Edition of Robot-Tasks : single loop, multi-rate
- Procedures : Semi-automatic generation of Esterel, user-level language (MaesTro, no longer maintained)
- GUI built on IlogViews
- Generation of C++ classes or C structures
- targets : VxWorks, Solaris, Linux, RTAI
- Verification : generic properties, Atg, observers
- Monitoring interface
- Maintained by Inria-RA staff

Current improvements (and dreams)

- Open-Source re-design (Kernel + plugins)
- Generic simulation templates
- Code distribution (including fault tolerance)
- Teleoperation (e.g. distributed vision based robot control)
- Flexible scheduling based on feed-back scheduling
- Graphical specification of procedures (Grafcet or StateCharts like)
- More friendly verification tools
- Discrete events controller synthesis
- ...

Model Driven Architecture emerging standards and tools