Fixed point aperiodic tilings

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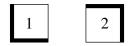
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Tiles and tilings

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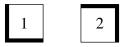
Tiles and tilings

Tile set:



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Tiling:

1	2	1	2	1
2	1	2	1	2
1	2	1	2	1
2	1	2	1	2
1	2	1	2	1

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C — a finite set of *colors* $\tau \subset C^4$ — a set of *tiles* (Wang tiles) tile $t = \langle t.left, t.right, t.up, t.down \rangle$ configurations: mappings $\mathbb{Z}^2 \to \tau$ *tilings*: configuration that satisfy matching rules C(i, j).right = C(i + 1, j).left C(i, j).up = C(i, j + 1).down

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A tiling C is *periodic* if it has some *period* T:

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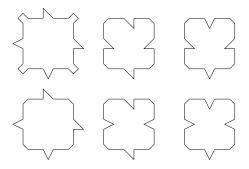
Aperiodic tile sets

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Theorem (Berger, 1966): there exists a tile set that has tilings but only aperiodic ones

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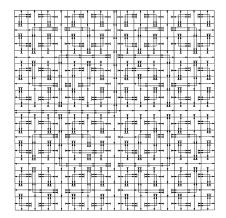
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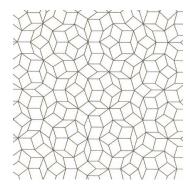
(Robinson tile set)

Tiling



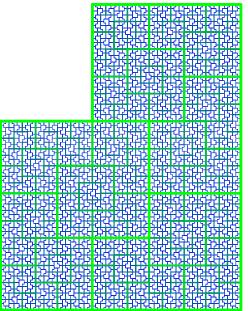
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Penrose tiling

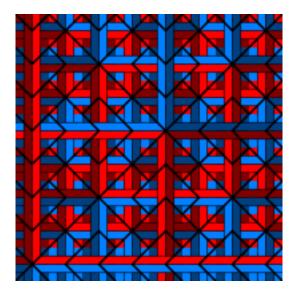


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Ammann tiling



Ollinger tiling



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 the question was asked by Hao Wang when he studied decision problems

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- Berger's construction became an important tool to prove undecidability of many algorithmic problems
- Aperiodic tiling can be constructed using self-referential argument widely used in logic and computation theory (Kleene's fixed point theorem)

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Self-similar tile sets

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Tile set τ is *self-similar* if it implements some set of macro-tiles ρ that is isomorphic to τ (Isomorphism: 1-1-correspondence that preserves matching rules)

Berger's theorem and self-similar tile sets

Berger theorem follows from two statements:

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B. There exists a self-similar tile set.

Let τ be a self-similar tile set with zoom factor M

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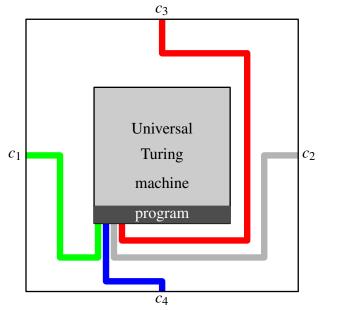
T/M is a multiple of M etc.

Self-referential tile set

- \blacktriangleright For a given tile set σ we construct a tile set τ that implements σ
- This gives a mapping $\sigma \to \tau(\sigma)$
- It remains to find a fixed point:

 $au(\sigma)$ is isomorphic to σ

The structure of a macro-tile that implements itself



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- robust aperiodic tile sets (isolated or sparse holes can be patched)
- simple proof of the unidecidability of the domino problem
- simple construction of a tile set that has only complex tilings
- tile set with any computable density