

Ergodic-type characterization of randomness

Laurent Bienvenu (LIAFA, Paris),

Adam Day (Victoria University, Wellington)

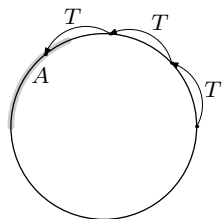
Ilya Mezhirov (Technical Univ., Kaiserslautern)

Alexander Shen (LIF, Marseille)¹

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¹on leave from IITP, Moscow

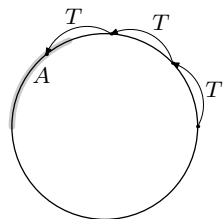
Ergodic theorem



X : space with measure μ

$T: X \rightarrow X$: measure preserving
 $x, T(x), T(T(x)), T^3(x), \dots$
how often in A ?

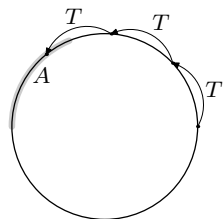
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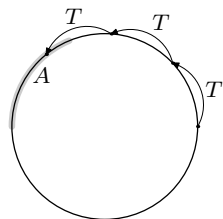


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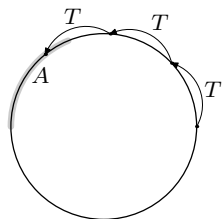
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answer: $\log_{10} 4 - \log_{10} 3$

Classical and algorithmic statements

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What we do: wider class of A s, but weaker statement:

$\mu(A) < 1 \Rightarrow$ for every random x at least one of $x, T(x), T^2(x), \dots$ is not in A

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\Leftrightarrow If $T^n(x) \in A$ for every n , then x is not Martin-Löf random

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effectively null set N : for every $\varepsilon > 0$ one can effectively generate a sequence of intervals that cover N and have total measure $< \varepsilon$

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Reformulation of Kucera's theorem: the set of all sequences x such that all tails of x are in A , is an effectively null set.

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(effective Kolmogorov 0-1-law)
- ▶ add some finite prefix to x to get $x' \notin A$

Each of these properties can be used as
characterization of randomness

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(In the proceedings T is required to be bijective; M. Hoyrup noted that it is not important.)

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Stronger claims in special cases: there are infinitely many shifts that move x outside A among any enumerable sequence of integers; statements about density of terms outside A

Application: Mijabe's result made easy

Theorem (Mijabe): let x^0 be a ML-random sequence, let x^1 be a ML-random sequence with oracle x^0 , let x^2 be a ML-random sequence with oracle x^0, x^1 etc. Then one can change finitely many terms in each x^i in such a way that x^0, x^1, \dots is a random element of $\Omega \times \Omega \times \dots$.

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Finite changes can be replaced by adding/deleting prefixes

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which is estimated using ergodic theorem for I and computability