



Oriented vertex and arc colorings of partial 2-trees

Pascal Ochem¹ Alexandre Pinlou²

*LaBRI, Université Bordeaux 1
351 Cours de la Libération
33405 Talence Cedex, France*

1 Introduction

We consider finite simple *oriented graphs*, that is digraphs with no opposite arcs. For an oriented graph G , we denote by $V(G)$ its *set of vertices* and by $A(G)$ its *set of arcs*. The number of vertices of G is the *order* of G . The *girth* of a graph G is the size of a smallest cycle in G . We denote by \mathcal{T}_g the class of partial 2-trees (also known as series-parallel graphs) with girth at least g .

The notion of oriented vertex-coloring was introduced by Courcelle [2] as follows: an *oriented k -vertex-coloring* of an oriented graph G is a mapping φ from $V(G)$ to a set of k colors such that (i) $\varphi(u) \neq \varphi(v)$ whenever $uv \in A(G)$ and (ii) $\varphi(v) \neq \varphi(x)$ whenever $uv, xy \in A(G)$ and $\varphi(u) = \varphi(y)$. The *oriented chromatic number* of G , denoted by $\chi_o(G)$, is defined as the smallest k such that G admits an oriented k -vertex-coloring. The oriented chromatic number $\chi_o(\mathcal{F})$ of a class of oriented graphs \mathcal{F} is defined as the maximum of $\chi_o(G)$ taken over all graphs G in \mathcal{F} .

Let G and H be two oriented graphs. A *homomorphism* from G to H is a mapping φ from $V(G)$ to $V(H)$ that preserves the arcs: $\varphi(u)\varphi(v) \in A(H)$ whenever

¹ Email: Pascal.Ochem@labri.fr

² Email: Alexandre.Pinlou@labri.fr

$uv \in A(G)$. An oriented k -vertex-coloring of an oriented graph G can be equivalently defined as a homomorphism ϕ from G to H , where H is an oriented graph of order k . The oriented chromatic number of G can then be viewed as the smallest order of an oriented graph H such that G admits a homomorphism to H . Links between colorings and homomorphisms are presented in more details in the monograph [3] by Hell and Nešetřil.

Oriented vertex-colorings have been studied by several authors in the last decade and the problem of bounding the oriented chromatic number has been investigated for various graph classes (see *e.g.* [1,8,9]).

Concerning partial 2-trees, Sopena proved [9] that their oriented chromatic number is at most 7 (this bound was shown to be tight). Pinlou and Sopena [8] obtained tight bounds for the oriented chromatic number of outerplanar graphs with given girth (outerplanar graphs form a strict subclass of partial 2-trees). Moreover, they proved that $\chi_o(\mathcal{T}_g) = 7$ for every g , $3 \leq g \leq 4$. In this paper, we complete the characterization of the oriented chromatic numbers of partial 2-trees with given girth:

Theorem 1.1

- (1) $\chi_o(\mathcal{T}_g) = 6$ for every girth g , $5 \leq g \leq 6$;
- (2) $\chi_o(\mathcal{T}_g) = 5$ for every girth g , $g \geq 7$;

One can define oriented arc-colorings of oriented graphs in a natural way by saying that, as in the undirected case, an *oriented arc-coloring* of an oriented graph G is an oriented vertex-coloring of its line digraph $LD(G)$ (recall that $LD(G)$ is given by $V(LD(G)) = A(G)$ and $ab \in A(LD(G))$ whenever $a = uv$ and $b = vw$). Therefore, an oriented arc-coloring ϕ of G must satisfy (i) $\phi(uv) \neq \phi(vw)$ whenever uv and vw are two consecutive arcs in G , and (ii) $\phi(vw) \neq \phi(xy)$ whenever $uv, vw, xy, yz \in A(G)$ with $\phi(uv) = \phi(yz)$. The *oriented chromatic index* of G , denoted by $\chi'_o(G)$, is defined as the smallest order of an oriented graph H such that $LD(G)$ admits a homomorphism to H . The oriented chromatic index $\chi'_o(\mathcal{F})$ of a class of oriented graphs \mathcal{F} is defined as the maximum of $\chi'_o(G)$ taken over all graphs G in \mathcal{F} .

The oriented chromatic index of oriented graphs was recently studied and several upper and lower bounds are known (see [6,7,8]).

Upper bounds for the oriented chromatic index can be easily derived from oriented chromatic number:

Claim 1.2 [6] For every oriented graph G , $\chi'_o(G) \leq \chi_o(G)$.

Our second result gives estimates of the oriented chromatic indexes of partial 2-trees with girth 4, 5 and 6, and a characterization for all other girths:

Theorem 1.3

- (1) $\chi'_o(\mathcal{T}_3) = 7;$
- (2) $6 \leq \chi'_o(\mathcal{T}_4) \leq 7;$
- (3) $5 \leq \chi'_o(\mathcal{T}_g) \leq 6$ for every girth $g, 5 \leq g \leq 6;$
- (4) $\chi'_o(\mathcal{T}_g) = 5$ for every girth $g, 7 \leq g \leq 17;$
- (5) $\chi'_o(\mathcal{T}_g) = 4$ for every girth $g, g \geq 18;$

In the rest of the paper, we will use the following notation. A vertex of degree k will be called a k -vertex. We denote by $\delta(G)$ the minimum degree of the graph G .

A k -path in a graph G is a path $P = [u, v_1, v_2, \dots, v_{k-1}, w]$ of length k (i.e. a path with k arcs) ; the vertices u and w are the endpoints of P . Note that a 1-path is an arc. A (k, d) -path is a k -path such that all internal vertices v_i have degree d .

A 2-vertex contraction is the contraction of an edge incident to a 2-vertex.

2 Sketches of proof

The proofs of Theorems 1.1 and 1.3 use some structural properties on partial 2-trees with given girth and on graph classes closed under 2-vertex contraction. These properties are given in the two following lemmas.

Lemma 2.1 *Let \mathcal{C} be a graph class closed under 2-vertex contraction such that every non-empty graph $G \in \mathcal{C}$ with girth at least g contains either a 1-vertex or a $(k, 2)$ -path, for some $k \geq 2$. Then, for every $n \geq 0$, every non-empty graph $G' \in \mathcal{C}$ with girth at least $g + n \lfloor \frac{g-1}{k-1} \rfloor$ contains either a 1-vertex or a $(k + n, 2)$ -path.*

For a graph G with girth at least g and a vertex $v \in V(G)$, we denote:
 $D_g^G(v) = |\{u \in V(G), d(u) \geq 3$ such that there exists a unique path of 2-vertices linking u and v or u and v are the endpoints of at least a $(\lfloor \frac{g}{2} \rfloor, 2)$ -path $\}|$.

Lemma 2.2 *Let G be a partial 2-tree with girth g such that $\delta(G) \geq 2$. Then, either there exists a $(\lfloor \frac{g}{2} \rfloor + 1, 2)$ -path, or there exists a ≥ 3 -vertex v such that $D_g^G(v) \leq 2$.*

Note that this lemma generalizes Lemma 2 p. 305 of Lih et al. [4] which characterizes partial 2-trees with girth 3.

Upper bounds

Thanks to the above lemmas, the upper bounds of Theorems 1.1 and 1.3 are obtained by showing that the considered partial 2-trees admit a homomorphism to one of the tournaments T_4, T_5, T_6 , and T_7 depicted on Fig. 1.

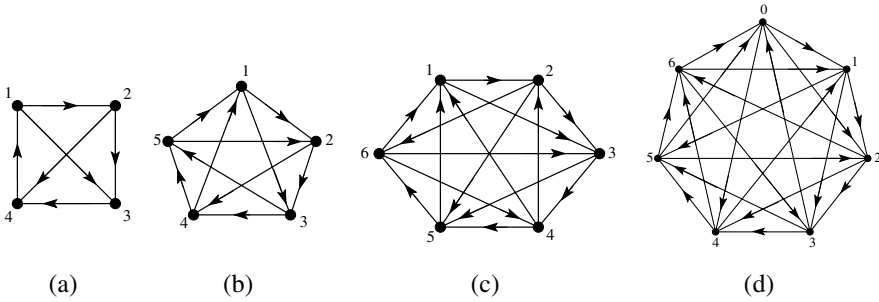


Fig. 1. The four target tournaments.

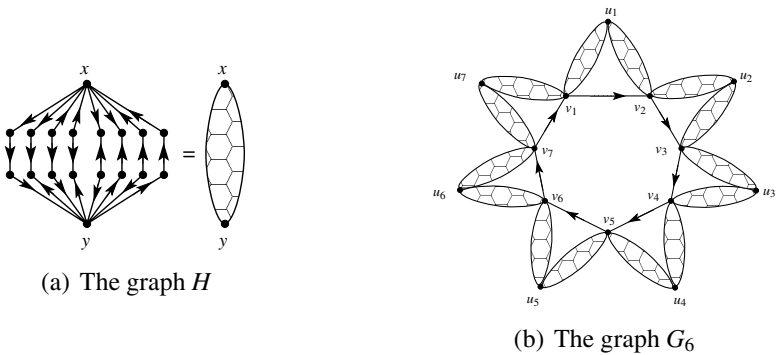


Fig. 2. An oriented partial 2-tree with girth 6 and oriented chromatic number 6.

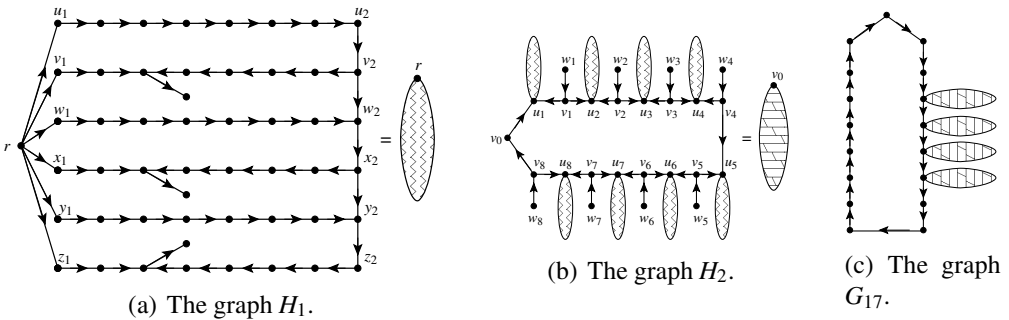


Fig. 3. An oriented partial 2-tree with girth 17 and oriented chromatic index 5.

Lower bounds

Finally, to get the lower bounds of Theorems 1.1 and 1.3, we construct partial 2-trees with the required girth which need the specified number of colors. More fully :

- The graph G_6 depicted in Fig. 2(b) is a partial 2-tree with girth 6 such that

$\chi_o(G_6) = 6$. Therefore, $\chi_o(\mathcal{T}_g) \geq 6$ for every $g \leq 6$.

- Nešetřil *et al.* [5] constructed for every $g \geq 3$, an oriented outerplanar graph with girth g which has oriented chromatic number 5. Therefore, $\chi_o(\mathcal{T}_g) \geq 5$ for every $g \geq 7$.
- The first three assumptions of Theorem 1.3 directly follow from Claim 1.2, Theorem 1.1(1) and some results of Pinlou and Sopena [8], namely $\chi_o(\mathcal{T}_3) = 7$, $\chi'_o(\mathcal{O}_4) = 6$, and $\chi'_o(\mathcal{O}_6) = 5$.
- The graph G_{17} depicted in Fig. 3(c) is a partial 2-tree with girth 17 such that $\chi'_o(G_{17}) = 5$. Therefore, $\chi'_o(\mathcal{T}_g) \geq 5$ for every $g \leq 17$.
- It not difficult to check that, for every $g \geq 3$, the partial 2-tree G obtained from two vertex-disjoint circuits, the first one of size g and the second one of size $k \geq g$ with $k \not\equiv 0 \pmod{3}$ has girth g and $\chi'_o(G) = 4$. Therefore $\chi'_o(\mathcal{T}_g) \geq 4$ for every $g \geq 18$.

References

- [1] O. V. Borodin, A. V. Kostochka, J. Nešetřil, A. Raspaud and É. Sopena, *On the maximum average degree and the oriented chromatic number of a graph*, *Discrete Math.* **206** (1999), pp. 77–89.
- [2] B. Courcelle, *The monadic second order-logic of graphs VI : on several representations of graphs by relational structures*, *Discrete Appl. Math.* **54** (1994), pp. 117–149.
- [3] P. Hell and J. Nešetřil, “Graphs and homomorphisms,” *Oxford Lecture Series in Mathematics and its Applications* **28**, Oxford University Press, 2004.
- [4] K. W. Lih, W. F. Wang and X. Zhu, *Coloring the square of K_4 -minor free graph*, *Discrete Math.* **269** (2003), pp. 303–309.
- [5] J. Nešetřil, A. Raspaud and É. Sopena, *Colorings and girth of oriented planar graphs*, *Discrete Math.* **165-166** (1997), pp. 519–530.
- [6] P. Ochem, A. Pinlou and É. Sopena, *On the oriented chromatic index of oriented graphs*, *Research Report RR-1390-06*, LaBRI, Université Bordeaux 1 (2006).
- [7] A. Pinlou, *On oriented arc-coloring of subcubic graphs*, *Elect. J. Comb.* **13** (2006).
- [8] A. Pinlou and É. Sopena, *Oriented vertex and arc colorings of outerplanar graphs*, *Inform. Process. Lett.* **100** (2006), pp. 97–104.
- [9] É. Sopena, *The chromatic number of oriented graphs*, *J. Graph Theory* **25** (1997), pp. 191–205.