

Random number generator: testing and whitening

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ESCAPE / LIRMM

invitation to an internship (2020)

What is a random sequence of digits?



● 000

- 000
- 01
- 1101111011100001101101010110001111101010
- 1100100100001111110110101010001000100001
(btw, the last sequence is a binary expansion of π)

Can you guess which sequence here is truly random?

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$x_1 \dots x_n$ is random means $C(x_1 \dots x_n) \approx n$

$x_1 \dots x_n$ is non-random means $C(x_1 \dots x_n) \ll n$

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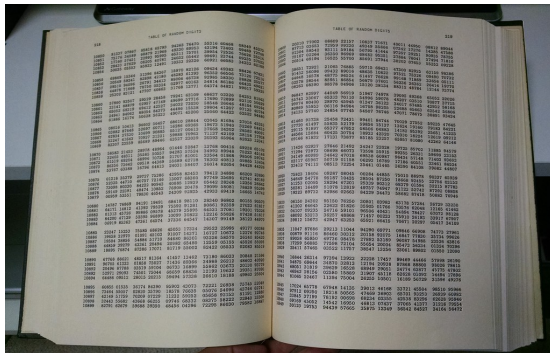
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Big practical question: where do we get truly random bits?



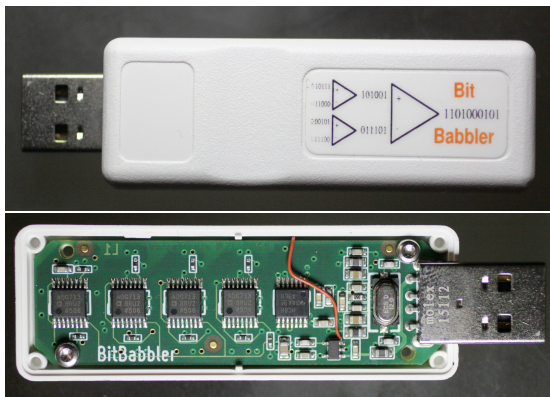
2

TABLE OF RANDOM DIGITS

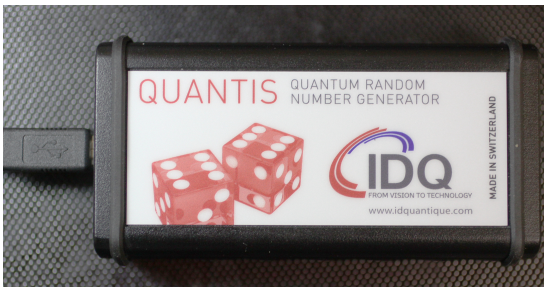
| | | | | | | | | | | |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| 00050 | 09188 | 20097 | 32825 | 39527 | 04220 | 86304 | 83389 | 87374 | 64278 | 58044 |
| 00051 | 90045 | 85497 | 51981 | 50654 | 94938 | 81997 | 91870 | 76150 | 68476 | 64659 |
| 00052 | 73189 | 50207 | 47677 | 26269 | 62290 | 64644 | 27124 | 67018 | 41361 | 82760 |
| 00053 | 75768 | 76490 | 20971 | 87749 | 90429 | 12272 | 95375 | 05871 | 93823 | 43178 |
| 00054 | 54016 | 44056 | 66281 | 31003 | 00682 | 27398 | 20714 | 53295 | 07706 | 17813 |

Rand Corporation, *A Million Random Digits with 100,000 Normal Deviates* (1955)

[random digits kindly generated for us in 1955]



[random digits from a noise in an electric circuit]

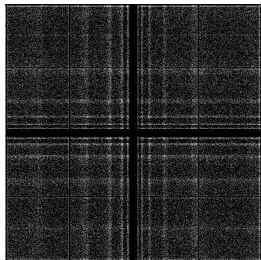


[random bits from quantum phenomena]

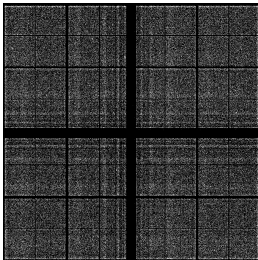
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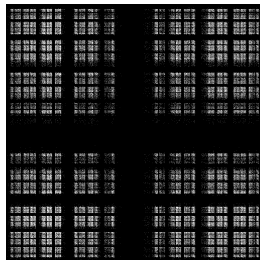
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100 kHz

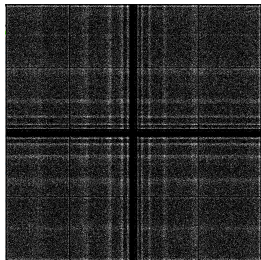


default rate 2.5 MHz

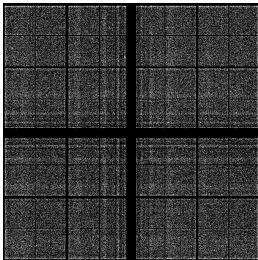


5 MHz

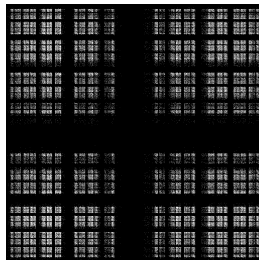
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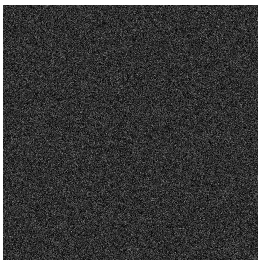
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successful whitening: XOR of 3 (apparently biased) data flows looks pretty random

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You cannot trust blindly the standard implementations of randomness tests (mathematically unsound tests, errors in the code).

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- apply (certainly non-perfect) pseudo-random generators to produce **useful** random objects (e.g., error correcting codes)

Internship proposal:

Generation of random bits is a classical problem known in the context of pseudo-random generators and also in connection with of truly random physical processes (there exist electronic devices that produce random bits using an unpredictable physical noise or intrinsically nondeterministic quantum phenomena). However, the quality of physical generators of random bits remains badly founded and poorly tested. The first objective of this project is an experimental study of the validity and quality of several physical **random numbers generators**.

When we talk about the quality of random or pseudo-random generators, we have to use **randomness tests**. The second objective of the project is an inventory and revision of statistical tests for random and pseudo-random generators. We suggest to improve the quality of statistical tests and develop new techniques of “whitening” that improves the quality of non-ideal sources of random bits. Another axis of the project is a conversion of various probabilistic proofs into unconventional randomness tests.

Prerequisites: Basic knowledge of probability and statistics, and **solid programming skills**. The main tools in the project are pretty standard: C / gcc / Linux. The project requires not only writing your own code but also reading and maintaining the code that already exists.